

06-14-00

A

UTILITY PATENT APPLICATION TRANSMITTAL
(Only for new nonprovisional applications under 37 CFR 1.53(b))

06/12/00

Docket No. : 39299/JEC/X2
Inventor(s) : Mahadevan Iyer, Lavanya Apsani and Pankaj Malviya
Title : OBJECT MODEL FOR NETWORK POLICY MANAGEMENT
Express Mail Label No. : EL386999355US

PC 09/592165
06/12/00

ADDRESS TO: Assistant Commissioner for Patents

Box Patent Application
Washington, D.C. 20231

Date: June 12, 2000

1. X FEE TRANSMITTAL FORM (Submit an original, and a duplicate for fee processing).

2. IF A CONTINUING APPLICATION

 This application is a of patent application No. .

Prior application information: Examiner ; Group Art Unit:

X This application claims priority pursuant to 35 U.S.C. §119(e) and 37 CFR §1.78(a)(4), to provisional Application Nos. 60/138,849, 60/138,850, 60/139,033, 60/139,034 60/139,035, 60/139,036, 60/139,038, 60/139,042, 60/139,043, 60/139,044, 60/139,047, 60/139,048, 60/139,049, 60/139,052, 60/139,053, all filed on June 10, 1999, and U.S. provisional application 60/139,076, filed on June 11, 1999

3. APPLICATION COMPRISED OF

Specification

54 Specification, claims and Abstract (total pages)

Drawings

30 Sheets of drawing(s) (FIGS. 1 to 30)

Declaration and Power of Attorney

 Newly executed
X No executed declaration
 Copy from a prior application (37 CFR 1.63(d))(for continuation and divisional)

4. Microfiche Computer Program (Appendix)

5. Nucleotide and/or Amino Acid Sequence Submission (if applicable, all necessary)

 Computer Readable Copy
 Paper Copy (identical to computer copy)
 Statement verifying identity of above copies

6. ALSO ENCLOSED ARE

 Preliminary Amendment
 A Petition for Extension of Time for the parent application and the required fee are enclosed as separate papers

UTILITY PATENT APPLICATION TRANSMITTAL
(Only for new nonprovisional applications under 37 CFR 1.53(b))

Docket No.: 39299/JEC/X2

- Small Entity Statement(s)
 - Statement filed in parent application, status still proper and desired
 - Copy of Statement filed in provisional application, status still proper and desired
- An Assignment of the invention with the Recordation Cover Sheet and the recordation fee are enclosed as separate papers
- This application is owned by pursuant to an Assignment recorded at Reel , Frame Information Disclosure Statement (IDS)/PTO-1449
 - Copies of IDS Citations
 - Certified copy of Priority Document(s) (*if foreign priority is claimed*)
 - English Translation Document (*if applicable*)
- Return Receipt Postcard (MPEP 503) (should be specifically itemized).
- Other

7. CORRESPONDENCE ADDRESS

CHRISTIE, PARKER & HALE, LLP, P.O. BOX 7068, PASADENA, CA 91109-7068

Respectfully submitted,

CHRISTIE, PARKER & HALE, LLP

By Josephine Chang
Josephine E. Chang
Reg. No. 46,083
626/795-9900

JEC/daa

OBJECT MODEL FOR NETWORK POLICY MANAGEMENT**CROSS-REFERENCE TO RELATED APPLICATIONS**

5 This application claims the benefit of U.S. provisional applications 60/138,849, 60/138,850, 60/139,033, 60/139,034 60/139,035, 60/139,036, 60/139,038, 60/139,042, 60/139,043, 60/139,044, 60/139,047, 60/139,048, 60/139,049, 60/139,052, 60/139,053, all filed on June 10, 1999, and U.S. provisional 10 application 60/139,076, filed on June 11, 1999, the contents of all of which are incorporated herein by reference.

FIELD OF THE INVENTION

15 The present invention relates to computer networks, and more particularly, to devices and methods for specifying and distributing policy management information effectively to remote private networks across the Internet.

BACKGROUND OF THE INVENTION

20 The growth and proliferation of computers and computer networks allow businesses to efficiently communicate with their own components as well as with their business partners, customers, and suppliers. However, the flexibility and efficiencies provided by such computers and computer networks 25 come with increasing risks, including security breaches from outside the corporation, accidental release of vital information from within it, and inappropriate use of the LAN, WAN, Internet, or extranet.

30 In managing the growth of computer networks as well as addressing the various security issues, network managers often turn to network policy management services such as firewall protection, Network Address Translation, spam email filtering, DNS caching, Web caching, virtual private network (VPN)

5 organization and security, and URL blocking for keeping network users from accessing certain Web sites through use of the organization's ISP. Each policy management service, however, generally requires a separate device that needs to be configured, managed, and monitored. Furthermore, as an organization grows
10 and spreads across multiple locations, the devices maintained also multiplies, multiplying the associated expenditures and efforts to configure, manage, and monitor the devices.

The solution to this problem is not as simple as just integrating multiple network policy management functions into a
15 single device at each location and allowing each location to share its policy information with other locations. In fact, there are many obstacles and challenges in adopting such an approach. One of these challenges is devising a scheme for specifying and distributing policy management information
20 effectively across the entire organization.

Accordingly, there remains a need in the art for an object model that allows policy management information to be specified and distributed effectively to remote private networks across the Internet.

25

SUMMARY OF THE INVENTION

The present invention is directed to a unified policy management system where various policies, namely, the set of rules and instructions that determine the network's operation,
30 may be established and enforced from a single site. According to one embodiment of the invention, the system includes a first edge device associated with a first network having a first set of resources that is configured to manage the policies for the first network according to the policy settings stored in a first database. The system also includes a second edge device
35 associated with a second network having a second set of resources

5 that is configured to manage the policies for the second network according to the policy settings stored in a second database. The first and second edge devices act as policy enforcers for their respective networks.

10 The system further includes a central policy server defining the first and second policy settings and managing the first and second edge devices from a single location. Thus, a network administrator need not multiply his or her efforts and associated expenditures in configuring and managing the policy enforcers individually.

15 The central policy server is also associated with a central database storing configuration information of the first and second edge devices. The central database is organized according to a hierarchical object oriented structure for simplifying policy management.

20 According to one aspect of the invention, the central database and the first and second databases are Lightweight Directory Access Protocol (LDAP) databases all adhering to the hierarchical object oriented structure.

25 According to another aspect of the invention, the structure includes a plurality of resource objects and policy objects for defining the first and second policy settings. The resource objects preferably include devices, users, hosts, services, and time. Devices are the policy enforcers at the edge of a particular private network and are associated with a set of users 30 and a particular host. A host is a network (e.g. a LAN subnet) in an organization. Services reflect the various services (e.g. HTTP, TELNET, FTP) provided by the policy server. Time is another dimension in controlling access to the network resources.

35 The policy objects preferably include bandwidth, firewall, administration, and virtual private network grouping policies. The virtual private network grouping is referred to as a VPN

5 cloud and includes one or more sites, users, and rules. A site includes one or more networks behind an edge device. The rules are firewall rules providing access control over network traffic flowing through the virtual private network.

10 It should be appreciated, therefore, that the hierarchical object oriented structure of the present invention helps simplify policy management by allowing the various elements of the policy management system to be defined and organized in an intuitive and extensible fashion.

15

BRIEF DESCRIPTION OF THE DRAWINGS

20 These and other features, aspects and advantages of the present invention will be more fully understood when considered with respect to the following detailed description, appended claims and accompanying drawings wherein:

FIG. 1 is a schematic block diagram of an exemplary unified policy management system;

25 FIG. 2 illustrates the hierarchical object-oriented structure of policies stored for an organization in accordance with the principles of the invention;

FIG. 3 is a schematic block diagram of a policy server in the policy management system of FIG. 1;

FIG. 4 is a schematic diagram of a central management sub-module in the policy server of FIG. 3;

30 FIG. 5 is an exemplary flow diagram of a device registration process carried out by the central management sub-module of FIG. 4;

FIG. 6 is a screen illustration of an exemplary graphical user interface for registering a device;

5 FIG. 7 is a screen illustration of an exemplary global monitor user interface presenting device health and status information;

10 FIG. 8 is a screen illustration of an exemplary graphical user interface provided by a policy management sub-module in the policy server of FIG. 3;

FIG. 9 is a screen illustration of an exemplary graphical user interface for managing system devices;

15 FIG. 10 is a screen illustration of an exemplary graphical user interface for managing system hosts;

FIG. 11 is a screen illustration of an exemplary graphical user interface for managing system services;

FIG. 12 is a screen illustration of an exemplary graphical user interface for managing time groups;

20 FIG. 13 is a screen illustration of an exemplary graphical user interface displaying a plurality of VPN clouds;

FIG. 14 is a screen illustration of an exemplary graphical user interface for adding a new firewall policy;

25 FIG. 15 is a schematic functional block diagram of policy enforcers updating their respective VPN membership information;

FIG. 16 is a block diagram of components in a self-extracting executable for downloading by a remote VPN client;

FIG. 17 is a functional block diagram for downloading the self-extracting executable of FIG. 16;

30 FIG. 18 is a schematic block diagram of a policy enforcer in the policy management system of FIG. 1;

FIG. 19 is a more detailed schematic block diagram of a policy engine in the policy enforcer of FIG. 18;

FIG. 20 is a more detailed schematic block diagram of a protocol classification engine of the policy enforcer of FIG. 18;

RECORDED-SEARCHED-COPYED-INDEXED-FILED

5 FIG. 21 is a more detailed schematic block diagram of an Internet protocol security engine in the policy enforcer of FIG. 18;

FIG. 22 is a schematic layout diagram of a common log format according to one embodiment of the invention;

10 FIG. 23 is a block diagram of an LDAP tree structure according to one embodiment of the invention;

FIG. 24 is a more detailed block diagram of a branch of the LDAP tree of FIG. 23;

15 FIG. 25 is a flow diagram for logging and propagating LDAP changes to policy enforcers;

FIG. 26 is a schematic block diagram of a high availability system including a primary unit and a backup unit;

FIG. 27 is a flow diagram of an exemplary status discovery process conducted by a high availability unit;

20 FIG. 28 is a flow diagram of a process for maintaining configuration information synchronized in the primary and backup units of FIG. 26;

FIG. 29 is an exemplary flow diagram of updating the primary and backup units of FIG. 26 when they are both functional; and

25 FIG. 30 is an exemplary flow diagram of updating the primary and backup units FIG. 26 when the primary is not functional.

DETAILED DESCRIPTION OF THE INVENTION

I. UNIFIED POLICY MANAGEMENT SYSTEM ARCHITECTURE

30 FIG. 1 is a schematic block diagram of an exemplary unified policy management system according to one embodiment of the invention. As illustrated in FIG. 1, private local networks 102, 104, and 106 are all coupled to a public network such as the Internet 108 via respective routers (generally identified at 110) and Internet Service Providers (ISPs) (not shown). Also coupled to the public Internet 108 via the ISPs are web surfers 112,

5 dial-up network users 114, servers providing unauthorized web sites 116, email spammers 118 sending out unsolicited junk email, and remote VPN clients 140 seeking access to the private local networks 102.

10 According to one example, local network 102 connects users and resources, such as workstations, servers, printers, and the like, at a first location of the organization, such as the organization's headquarters, and local network 104 connects users and resources at a second location of the organization, such as a branch office. Furthermore, local network 106 connects users 15 and resources of a customer of the organization requiring special access to the organization's users and resources. Authorized dial-up network users 114 of the organization are respectively situated at remote locations from the first and second local networks, and also require special access to the organization's 20 users and resources. Furthermore, web surfers 112 communicate with the organization's web server 120 over the public Internet 108 and access the organization's web site.

Local network 102 includes a policy server 122 for defining and managing network services and policies for the organization. 25 The network policies are a set of rules and instructions that determine the network's operation, such as firewall, VPN, bandwidth, and administration policies. The firewall policies decide the network traffic that is to be allowed to flow from the public Internet 108 into the local networks 102, 104, and the 30 traffic that is to be blocked. The bandwidth policies determine the kind of bandwidth that is to be allocated to the traffic flowing through the local networks. The VPN policies determine the rules for implementing multiple site connectivity across the local networks. The administration policies decide the users 35 that have access to administrative functions, the type of administrative functions allocated to these users, and the policy

5 enforcers 124, 126 on which these users may exercise such administrative functions. The firewall, VPN, bandwidth, and administration policies for the entire organization are preferably stored in a policy server database 130 maintained by the policy server 122.

10 Each local network 102, 104 also includes an edge device, referred to as a policy enforcer 124, 126, for controlling access to the network. Each policy enforcer 124, 126 manages the network policies and services for the users and resources of their respective local networks 102, 104, as permitted by the 15 policy server 122. Respective portions of the policy server database 130 are copied to the policy enforcer databases 132, 134 for allowing the policy enforcers to manage the network policies and services for the local networks 102, 104.

20 According to one embodiment of the invention, the policy server 122 and policy enforcers 124, 126 may be implemented in a similar fashion as the FORT KNOX series of policy routers made by Alcatel Internetworking, Inc., of Milpitas, California.

25 **II. OBJECT MODEL FOR NETWORK POLICY MANAGEMENT**

According to one embodiment of the invention, the policy server database 130 and policy enforcer databases 132, 134 are LDAP databases adhering to a unified hierarchical object oriented structure. The LDAP directory service model is based on entries 30 where each entry is a collection of attributes referenced by a distinguished name (DN). Each of the attributes includes a type and one or more values. The type is typically a mnemonic string, such as "o" for organization, "c" for country, or "mail" for email address. The values depend on the type of attribute. For 35 example, a "mail" attribute may contain the value "babs@umich.edu." A "jpegPhoto" attribute may contain a photograph in binary JPEG/JFIF format. Additional details of the

5 LDAP directory service model are defined in RFC 1777 "The
Lightweight Directory Access Protocol" (W. Yeong, T. Howes, and
Kille, Network Working Group, March 1995) and "LDAP Programming:
Directory-enabled Applications with Lightweight Directory Access
Protocol" (T. Howes, and M. Smith, Macmillan Technical
10 Publishing, 1997), incorporated herein by reference.

The entries in the LDAP database are preferably arranged in
a hierarchical tree-like structure reflecting political,
geographic, and/or organizational boundaries. Entries
representing countries appear at the top of the tree. Below them
15 are entries representing states or national organizations. Below
the states or national organizations may be entries representing
people, organization units, printers, documents, and the like.

FIG. 2 is a schematic layout diagram of a unified
hierarchical object oriented structure adhered by the policy
20 server database 130 according to one embodiment of the invention.
The policy enforcer databases 132, 134 adhere to a similar
structure except for a few differences. For example, the policy
enforcer databases preferably do not contain a policy server
domain object 201 and related policy server objects, nor a policy
25 domain object 240.

As illustrated in FIG. 2, each object in the structure is
preferably stored as an LDAP entry. At the top of the hierarchy
is the policy server domain object 201 including various policy
server resources and a plurality of policy domains objects
30 (generally referenced at 204). Each policy domain object 240 is
a grouping of policy enforcers that share common policies. Each
policy domain object 240 includes a resource root object 200 and
a group root object 202. All policy management functions are
preferably implemented in terms of the resource objects which
35 include devices 204, users 206, hosts 208, services 210, and time
220. Thus, a firewall policy may be defined by simply assigning

5 the particular devices, users, hosts, services, and time
applicable to the policy. The devices, users, hosts, and
services are preferably organized in groups 212, 214, 216, and
218, respectively, having a group name, description, and member
information for a more intuitive way of addressing and organizing
10 the resources.

Users 206 are preferably associated with a user domain
providing a secure and efficient means of authenticating the
user. Each user domain has a single policy enforcer who is
authorized to authenticate the user. Thus, user domains ensure
15 that the authenticating agent is generally located in the same
local network as the user. This helps eliminate the cost of
network dependency or network latency during the user
authentication process. It should be noted, however, that users
may also constitute authorized dial-up users 114 and users from
20 the customer network 106. These users contact a remote
authenticating agent which proxies the authentication back to the
appropriate policy enforcer.

Hosts 208 are the various networks present in an
organization. For instance, a particular LAN subnet may be
25 specified as a host in the system. Hosts 208 are preferably
organized based on their physical locations within the
organization. A host's physical location is identified by the
device (policy enforcer) 204 associated with the host.

Services 210 reflect the various services provided by the
30 policy server 122. Such services include, for example,
multimedia streaming/conferencing, information retrieval,
security and authentication, database applications, mail
applications, routing applications, standard communication
protocols, and the like. Attributes associated with each service
35 preferably include a service name, description, type (e.g. HTTP,

5 HTTPS, FTP, TELNET, SMTP, Real Networks, and the like), and
group.

10 Devices 204 are the policy enforcers 124, 126 at the edge
of a particular local network. Each device/policy enforcer
preferably includes users 206 and a host/network 208 that is
managed by the policy enforcer.

Time 220 is another dimension in controlling access to the
network resources. Various time objects covering a range of
times may be created and used in creating the firewall policies.

15 Similar to resources, network policies are also preferably
defined in terms of objects for a more efficient and intuitive
definition of the policies. Policies are defined by the
administrators and implemented by the policy enforcers 124, 126
on the network traffic flowing between the public Internet 108
and the local networks 102 and 104.

20 According to one embodiment of the invention, a policy
object 222 includes a bandwidth policy 224, firewall policy 226,
administration policy 228, and VPN policy 230. The VPN policy
230 defines a security policy for the member networks and
includes one or more VPN clouds 232. Each VPN cloud 232 is an
individual VPN or a group of VPNs defining a security policy
group which includes a list of sites 234 and users 236 who can
communicate with each other. A site is preferably a set of
hosts/networks physically located behind one of the policy
enforcers 124, 126. In other words, a site is a definition of
30 a network which includes the policy enforcer that is associated
with it. The policy enforcers for the sites act as VPN tunnel
endpoints once the hosts under the sites start communicating.
These communications are governed by a set of rules 238
configured for each VPN cloud. The rules 238 may govern, among
35 other things, VPN access permissions and security features such

5 as the level of encryption and authentication used for the connectivity at the network layer.

The object oriented structure of FIG. 2 thus allows the network administrators to define policies in an intuitive and extensible fashion. Such policies may be defined by simply 10 associating resources to the policies. This allows for a policy-centric management model where the administrator is given the impression that a single logical server provides the firewall, bandwidth management, and VPN services across the enterprise.

The fact that the policy is enforced on individual policy 15 enforcers in different locations is transparent to the administrator.

III. POLICY-BASED NETWORK ARCHITECTURE

FIG. 3 is a more detailed schematic block diagram of the 20 policy server 122 according to one embodiment of the invention. The policy server 122 preferably includes a management module 302 that allows centralized control over the policy enforcers 124, 126 from a single console. The policy server 122 further includes a log collecting and archiving module 304 and a policy 25 server reports module 316. The log collecting and archiving module 304 collects information about the status and usage of resources from the policy enforcers 124, 126 as well as from the management module 302, and stores them in an archive database 318. The policy server reports module 316 uses the collected 30 logs and archives to generate reports in an organized report format.

Referring again to the management module 302, the management 35 module 302 preferably includes four sub-modules aiding in the centralized control, namely, a centralized management sub-module 306, policy management sub-module 308, secure role-based

5 management sub-module 310, and multiple site connectivity management sub-module 312.

10 The centralized management sub-module 306 enables a network administrator to install and manage individual policy enforcers from a central location. The network administrator preferably uses a web-based graphical user interface to define the policy enforcer's network configuration and monitor various aspects of the device, such as device health, device alarms, VPN connection status, and the like.

15 The policy management sub-module 308 provides the network administrator with the ability to create policies that span multiple functional aspects of the policy enforcer (e.g. firewall, bandwidth management, and virtual private networks), multiple resources (e.g. users, hosts, services and time), and multiple policy enforcers.

20 The secure role-based management sub-module 310 provides role-based management to enable administrators to delegate administrative responsibilities to other administrators. This sub-module preferably provides for maximum security when it comes to accessing the management functions.

25 The multiple site connectivity management sub-module 312 allows the network administrator to set-up secure communication channels between two or more remote sites. In doing so, this sub-module leverages the centralized management sub-module 306, policy management sub-module 308, dynamic routing capabilities 30 of the policy enforcers 124, 126, and the management infrastructure to provide virtual private networks across the enterprise with fine grained access control.

35 FIG. 4 is a more detailed schematic diagram of the central policy management sub-module 306 according to one embodiment of the invention. The sub-module includes a policy server installation wizard 404 providing an interactive user interface

5 to aid the installation of the policy server 122. In this regard, the network administrator has access to a personal computer connected to a LAN port of the policy server 122 via a cross over cable, hub, or the like. The network administrator connects to the policy server 122 by preferably typing-in a URL 10 of the policy server 122 into a standard Internet browser such as Microsoft Internet Explorer. The URL is preferably of the form of "http://<ipaddress>:88/index.html" where <ipaddress> is the IP address that is to be assigned to the policy server. The IP address is automatically assigned to the policy server when 15 the browser attempts to contact the address. When the administrator's personal computer sends an address resolution protocol request for the IP address, the policy server detects that a packet directed to port 88 is not claimed, and assumes the IP address.

20 Once connected, the policy server installation wizard 404 invokes the interactive user interface to assist the administrator in setting up the policy server 122. Among other things, the policy server installation wizard 404 prompts the administrator to specify a server name, server IP address, and router IP address. Furthermore, the policy server installation 25 wizard 404 prompts the administrator to select one of various default policies for creating default firewall, VPN, bandwidth, and administrator policies. These policies are then replicated on each new policy enforcer registering with the policy server 30 122.

The centralized management sub-module 306 further includes a policy enforcer installation wizard 406 providing an interactive user interface to aid the installation of the policy enforcers 124, 126. As with the installation of the policy 35 server 122, the access to the wizard 406 is preferably web-based using the network administrator's personal computer.

5 Once connected, the policy enforcer installation wizard 406
invokes the interactive user interface to assist the network
administrator in setting up a particular policy enforcer 124,
126. Among other things, the policy enforcer installation wizard
10 464 prompts the administrator to specify the policy server IP
address, policy enforcer IP address, and router IP address. The
policy enforcer then registers with the policy server 122 by
invoking a URL on the policy server with basic bootstrap
information of its own. The registration of the policy enforcer
allows the initialization of the policy enforcer's database 132,
15 134 with the configuration information, as well as the monitoring
of the policy enforcer's status and health by the policy server
122.

Prior to registering the policy enforcer with the policy
server 122, the network administrator preferably pre-registers
20 the policy enforcer on the policy server. Such pre-registering
allows the creation of a placeholder node on the policy server
for the policy enforcer data for when the policy enforcer does
in fact register. In this regard, the centralized management
sub-module 306 includes a configuration interface 410 allowing
25 the pre-registration of a new policy enforcer.

FIG. 5 is an exemplary flow diagram of a policy enforcer
pre-registration and registration process according to one
embodiment of the invention. In step 401, the policy enforcer
is connected to the network and installed at its actual physical
30 location using the above-described policy enforcer installation
wizard 406. The network administrator, possessing the new
device's serial number, pre-registers the policy enforcer by
adding the new policy enforcer to a device group in step 403.
In this regard, the configuration interface 410 invokes an
35 interactive graphical interface, such as the one illustrated in
FIG. 6, allowing the network administrator to enter a device name

5 415, serial number 417, and location information 419, and further
allowing the administrator to select a device group 421 to which
the new policy enforcer is to belong. Actuation of an apply
button 423 causes the new policy enforcer, in step 405, to
contact the policy server 122 by preferably invoking a URL on the
10 policy server. Once the policy server has been contacted, the
new policy enforcer transmits its registration packet to the
policy server. The registration packet includes at least a
serial number of the new policy enforcer, as well as the IP
addresses of the LAN, WAN, and DMS on the policy enforcer. In
15 step 407, the centralized management sub-module 306 compares the
serial number of the new policy enforcer with the list of policy
enforcers pre-registered with the policy server 122. If a match
is found, the policy server 122 proceeds with the registration
process by packaging, in step 409, the settings selected for the
20 policy enforcer during its installation process, preferably into
an LDAP Data Interchange Format (ldif) file. In step 411, the
file is transmitted to the policy enforcer, preferably over an
HTTPS channel, by invoking a common gateway interface (CGI) on
the policy enforcer. The policy enforcer then uses the file to
25 initialize its configuration database, such as database 132, 134,
in step 413.

Referring again to FIG. 4, the centralized management sub-
module 306 also includes a global monitor user interface 402 and
a data collector program 412, respectively displaying and
30 collecting the health and status of all the policy enforcers
managed by the policy server 122. The data collector program 412
receives health and status information from each of the up-and-
running policy enforcers it manages, and passes the relevant
information to the global monitor user interface. A health agent
35 running as a daemon in each of the policy enforcers being
monitored periodically collects data from the device and analyzes

5 its health status. The collected data is then transferred to the policy server 122 when requested by the data collector program 412.

10 FIG. 7 is a screen illustration of an exemplary global monitor user interface 402 presenting various types of health and status information. Such information may relate to the health of the device, such as system load 712 and network usage information 714. The information may also relate to current alarms 716 on the device including alarm name, type, description, and the like. The information may further relate to current VPN 15 connections 718 including connection type, source/destination, duration, and VPN traffic volume.

20 Referring again to FIG. 3, the policy management sub-module 308 allows for policy management of the policy enforcers 124, 126. As discussed above, all policy management functions are implemented in terms of resource objects stored in the policy databases 130, 132, 134 including users, devices, hosts, services, and time. Preferably, all resources are associated with default policy settings selected by the administrator during the installation process. The network administrator views, adds, 25 and modifies the policies centrally via a graphical user interface provided by the policy management sub-module 308. This allows for a policy-centric management model where the administrator is given the impression that a single logical server provides the firewall, bandwidth management, and VPN 30 services across the enterprise. The fact that the policy is enforced on individual policy enforcers in different locations is transparent to the administrator.

35 FIG. 8 is a screen illustration of an exemplary graphical user interface provided by the policy management sub-module 308. The interface includes a resource palette 718 including a list of resource tabs including a users tab 718a, devices tab 718b,

5 hosts tab 718c, services tab 718d, and time tab 718e. The
resource palette allows the administrator to add and modify
resource definitions from a single console.

10 Selection of the users tab 718a causes a display of the user
groups 722 defined for the system. New users may be added to the
group by selecting a particular group and defining various
15 attributes of the user such as a login name, full name, policy
enforcer to which the user belongs, authentication scheme,
password, and the like.

15 Selection of the devices tab 718b causes a display of
various device management icons for managing the policy server
20 122 and the policy enforcers 124, 126 as is illustrated in FIG.
9. A policy server systems settings icon 750 allows the network
administrator to view and modify system settings like LAN,
WAN/DMS IP addresses of the policy server 122. A policy server
archive options icon 752 allows specification of reporting and
25 other database archive options at the policy server 122. A
global URL blocking icon 754 allows the administrator to specify
a list of unauthorized web sites 116 to be blocked by all the
policy enforcers 124, 126 of the system. Similarly, a global
spam list icon 756 allows the administrator to specify a list of
30 email addresses of spammers 118 to be blocked by all the policy
enforcers.

35 The administrator may view information on all the policy
enforcers 124, 126 by selecting icon 758. Information on a
specific policy enforcer may be viewed by selecting a specific
policy enforcer 760 under a particular device group 761. Such
information includes system settings information 762, URL
blocking information 764, spam list information 766, and the
like, that is specific to the selected policy enforcer. For
instance, selection of the policy enforcer's URL blocking
information 764 icon causes a display of various categories 768

5 of URLs that the network administrator may select to block for
the selected policy enforcer.

Selection of the hosts tab 718c causes a display of various hosts (networks) of the system as is illustrated in FIG. 10. A host is organized based on its physical location and is further
10 associated with a particular policy enforcer 124, 126. Hosts are associated with various attributes including a unique name 770, an IP address of the network 772, and a subnet mask 774. In addition, the administrator may specify whether the host is an external host 776 belonging to a network that is not administered
15 by the policy server 122. If the host is an external host, the administrator specifies an IP address 778 of the external device to which the host belongs. A device field 780 allows the administrator to enter the policy enforcer's name to which the host belongs. Each host is further associated with a particular
20 group 782 assigned by the administrator.

Selection of the services tab 718d causes a display of various service groups supported by the policy server 122 as is illustrated in FIG. 11. Such service groups include, for example, multimedia streaming/conferencing, information retrieval, security and authentication, mail applications, routing applications, database applications, standard communication protocols and the like. Users may also add new service groups as desired.

Each service is associated with a name 784, description 786, and service type 788 (e.g. HTTP, HTTPS, FTP, TELNET, SMTP, Real Networks, and the like). Furthermore, each service is associated with a service group 790. Based on the type of service, additional information may also be specified for the service. For instance, for an HTTP service, the administrator may specify
35 whether URL blocking 792 is to be enabled.

5 Selection of the time tab 718e causes a display of various time group icons 794 covering a range of times to be used in the firewall policies as is illustrated in FIG. 12. For instance, selection of a work time group icon allows the network administrator to set the days and times which are to be set as
10 working days and hours.

Referring again to FIG. 8, the interface also includes a policy canvas 720 including a list of policies available to the system. A policy definition is preferably an association of a set of resources that may be dragged from the resource palette
15 718 and dropped onto the policy canvas 720.

Selection of a firewall tab 720a causes a display of all the firewall policies defined for a particular policy domain including one or more policy enforcers. The network administrator decides the domain to which a policy enforcer is to belong during pre-registration of the policy enforcer. The interface allows the network administrator to view, add, and modify the various policies from the policy server 122 and effectuate the changes on the policy enforcers 124, 126 without the need to make such changes individually in each policy
20 enforcer.

According to one embodiment of the invention, each firewall policy includes a policy identifier (ID) attribute 724 for identifying a particular policy rule in the list of policies. An order number attribute 726 for the policy rule indicates the sequence in which the policy is to be applied. In this regard, the policy enforcer 124, 126 for the local network takes one rule
30 at a time, in sequence, compares it against the network traffic, and preferably applies the first rule that matches the network traffic.

35 Each firewall policy also includes a description attribute 728 for describing the firewall policy to be applied. For

5 instance, the description may indicate that the policy allows
spam blocking, URL blocking, VPN key management, and the like.
An action flag attribute 730 indicates whether traffic is to be
allowed or denied for the indicated policy. An active flag
attribute 732 indicates whether the policy has been activated
10 or de-activated. Thus, the network administrator may create a
policy and activate it at a later time. A policy that has been
de-activated preferably has no effect on the network traffic.

Each firewall policy further includes a user attribute 734,
source attribute 736, service attribute 738, destination
15 attribute (not shown), and time attribute (not shown). Each of
these attributes is preferably represented by a group name or a
resource name. The name acts as a pointer to an entry in the
group root object 202 or resource root object of the LDAP
database 130, 132, or 134.

20 Preferably, the user attribute 734 indicates the user groups
and users that are eligible for the policy. The source attribute
736 indicates a point of origination of the network traffic
associated with the user. The services attribute 738 indicates
the services to the allowed or denied by the policy. The
25 destination attribute indicates a specific LAN, WAN, DMS segment
or specific hosts where the specified services are to be allowed
or denied. For example, to configure SMTP pop services on a mail
server, the host may be the IP address where the mail server is
running, and the services specified is SMTP. The time attribute
30 indicates a time slot in which the policy is to be effective,

In addition to the above, each firewall policy also includes
an authentication attribute (not shown) indicating an
authentication scheme for the policy (e.g. none, LDAP, SecurID,
RADIUS, WinNT, or all).

35 FIG. 14 is a screen illustration of an exemplary graphical
user interface for adding a new firewall policy to the policy

5 domain upon actuation of an add button 725. Existing firewall policies may also be modified or deleted by actuation of a modify button 727 and a delete button 729, respectively.

As illustrated in FIG. 14, a new firewall policy may be defined by simply adding a description of the policy in a
10 description area 728a, selecting an action to be applied to the matching network traffic in an action box 730a, and indicating in an active area 732a whether the policy is to be active or inactive. Furthermore, the network administrator specifies the user, source, services, destination, and time resources in a user
15 area 734a, source area 736a, services area 738a, destination area 739a, and time area 741, respectively. The network administrator further selects an authentication scheme for the policy in an authentication area 743. Upon actuation of an OK button 745, appropriate entries of the policy server database's LDAP tree are
20 suitably changed to reflect the addition of the new policy. The change is also transmitted to the respective policy enforcers as is described in further detail below.

Referring again to FIG. 8, selection of the bandwidth tab 720c allows the display, addition, and modification of various
25 bandwidth policies determining the kind of bandwidth to be allocated to a traffic flowing through a particular policy enforcer. Different bandwidths may be specified for different users, hosts, and services.

Selection of the administration tab 720d allows the display, addition, and modification of various administrative policies allowing a head network administrator to delegate administrative responsibilities to other administrators. In this regard, the head network administrator specifies administration policies that determine which users have access to what functions, and for what
30 devices. Preferably the administration policies include similar attributes as the firewall rules except for the specification of
35

5 a role attribute. Extra administrative privileges may be afforded to certain users depending on their role.

IV. VIRTUAL PRIVATE NETWORK HAVING AUTOMATIC REACHABILITY UPDATING

10 Referring again to FIG. 3, the multi-site connectivity management module 312 allows the creation of dynamically routed VPNs where VPN membership lists are automatically created without statically configuring the membership information by the network administrator. Thus, once the administrator configures a VPN from one policy enforcer's LAN to another, routing protocols such as RIPv1 or RIPv2 running on the LAN interfaces learn about the networks reachable through their respective interfaces. These networks then become the VPN's members, and the policy enforcers 124, 126 on either side of the VPN create membership tables using the learned routes. The membership information is preferably exchanged between the policy enforcers 124, 126 through the LDAP databases 132, 134. Thus, the combined use of routing protocols and LDAP allows the creation of VPNs whose member lists are 20 dynamically compiled.

25

Referring again to FIG. 8, the network administrator configures VPN policies for multiple site connectivity using the resource palette 718 and policy canvas 720. Selection of the VPN tab 720b in the policy canvas 720 causes the display of a 30 collection of VPN clouds 270 already configured for the system as is illustrated in FIG. 13. As described above, a VPN cloud is an individual VPN or a group of VPNs for which a security policy may be defined. Each VPN cloud includes a list of sites under a sites node 234 and users under a users node 236, who can 35 communicate with each other. A site is a set of hosts that are physically behind one of the policy enforcers 124, 126. The

5 policy enforcers for the sites preferably act as VPN tunnel endpoints once the hosts under the sites start communicating.

10 The users in the VPN cloud are the users who may access the hosts associated with the sites 234. The users access the hosts as VPN clients using VPN client software installed in each user's personal computer as is described in further detail below.

15 Each VPN cloud 270 further includes a firewall rules node 276 including firewall rules to be applied all the connections in the cloud. The rules may govern, among other things, VPN access permissions, security features such as the level of encryption and authentication used for the connectivity at the network layer.

20 The hierarchical organization provided by the VPN clouds thus allows the network administrator to create fully meshed VPNs where every site within a VPN cloud has full connectivity with every other site. The network administrator need no longer manually configure each possible connection in the VPN, but only need to create a VPN cloud and specify the sites, users, and rules to be associated with the VPN. Each connection is then configured based on the configuration specified for the VPN cloud. The hierarchical organization thus facilitates the setup 25 of a VPN with a large number of sites.

30 The network administrator preferably adds a new VPN cloud by actuating an add button 280. In response, the policy server 122 automatically creates the sites node 272, users node 274, and rules node 276 under the VPN cloud. The administrator then specifies the sites and users in the VPN.

35 According to one embodiment of the invention, the rules node 276 initially includes a default VPN rule 278 corresponding to the policy settings selected by the network administrator during setup of the policy server 122. The default VPN rule 278 allows unrestricted access between the hosts in the VPN.

5 The administrator may implement the access control within
the VPN cloud by deleting the default rule 278 and adding
specific firewall rules to the VPN. Such firewall rules allow
the administrator to have fine grained access control over the
traffic that flows through the VPN, all within the realm of the
10 encrypted access provided by such VPN. The firewall rules are
applied to the cleartext packet after it is decrypted or before
it is encrypted.

15 According to one embodiment of the invention, the
administrator selects the default rule 278 to effectuate such
changes to the default rule. Selection of the default rule
invokes a graphical user interface similar to the one illustrated
in FIG. 8. The network administrator then fine tunes the access
to the VPN by defining the firewall rules applicable to the VPN.
The parameters in these firewall rules are preferably identical
20 to the general firewall rules illustrated in FIG. 8.

25 Once a VPN cloud is configured, VPN membership information
is dynamically created by the policy enforcers 124, 126 in the
VPN. In this regard, each VPN site includes a tag identifying
the hosts included in the site. At runtime, the policy enforcers
124, 126 for the respective sites associate IP addresses to the
tag identifying the hosts in each site. This allows the IP
addresses to be dynamically discovered without requiring static
configuration of the IP addresses.

30 After the creation of the membership tables, any changes in
the routing information is detected and notified to the member
policy enforcers using a publish/subscribe process. The actual
changes are retrieved by a policy enforcer by querying the LDAP
database on the particular network that corresponds to the
changed routing information.

35 FIG. 15 is a schematic functional block diagram of policy
enforcers 124, 126 at opposite ends of a VPN tunnel updating

5 their respective routing information. As illustrated in FIG. 15, each policy enforcer 124, 126 includes a gated module 252, 261 configured as a daemon to run one or more routing protocols for exchanging routes on the network. Such routing protocols may include RIPv1, RIPv2, OSPF, and the like.

10 When a network administrator wishes to add a new route to the private local network 102 connected to policy enforcer 124, the administrator submits, in step 241, the new route to a gated module 252 in the policy enforcer 124. This is typically done by configuring a downstream of the policy enforcer to have an 15 additional network. This information is then propagated by standard routing protocols to the gated module 252 of the policy enforcer 124. For example, the policy server 122 may publish the new route to the policy enforcer 124 with which the new route is to be associated. The route may be specified, for example, by 20 an LDAP statement such as "LAN_Group@PR1," which specifies a new route from a policy enforcer PR1 to a LAN named LAN_Group. The gated module 252, in step 242, writes the new route to a kernel 253 of the policy enforcer including a VPN driver 254 so that the policy enforcer 124 can properly direct appropriate messages 25 along the new route. Furthermore, the gated module 252, in step 243, writes the new route to its LDAP database 132.

The gated module 252 also provides, in step 244, the name of the new route to a distinguished name monitor (DNMonitor) 30 daemon 255 configured to listen for updates in the LDAP database 132. The DNMonitor in turn notifies, in steps 245a, 245b, a VPN daemon 256 and a policy deployment point (PDP) engine 257 of the change in the LDAP database 132. The PDP engine then updates the modules that enforce the policies, with the change.

The VPN daemon 256, in step 246, uses the route name to 35 access the LDAP database 132 to get the complete route information, a list of all VPNs to which the new route belongs,

5 and a list of all other policy routers connected to those VPNs. In step 247, the VPN daemon 256 proceeds to send the new route name to each of the other policy routers.

When policy router 126 receives a new route name from policy router 124, its network daemon 258, in step 248, accesses the 10 LDAP database 132 in the sending policy router 124 to obtain the complete new route information. If the new route belongs to more than one VPN and has different parameters for the different VPNs, routers on the different VPNs retrieve different information corresponding to the individual VPNs.

15 In step 249, the network daemon 258 writes the new route information obtained in its own LDAP database 134 and provides it to its own DNMonitor module. As in the sending policy router 124, the DNMonitor module 259 in the receiving policy router 126 provides the new route information to its PDP engine 260 for 20 updating its kernel 265 with the latest changes.

Although FIG. 15 has been described in connection with addition of a route to a policy enforcer and its associated VPNs, it should be readily apparent to those skilled in the art that 25 essentially the same techniques may be applied to deletion of a route (for example, if a network component becomes inoperative or incommunicative), or change of a route (the policy router may recognize that a route already exists in a different form and simply overwrite it). In this way, the VPN system or systems can dynamically maintain routing information between its policy 30 enforcers with minimal intervention by the system administrator.

V. VIRTUAL PRIVATE NETWORK HAVING AUTOMATIC UPDATING OF CLIENT REACHABILITY INFORMATION

35 Remote users communicate over the public Internet 108 with the other members of the VPN behind policy enforcers 124, 126, upon presenting appropriate credentials. These remote users

5 access the private networks as VPN clients 140 using a VPN client
software. According to one embodiment of the invention, the
system allows the remote user to download a self-extracting
executable which, upon execution, installs both the VPN client
software and VPN reachability information unique to the remote
10 user in the user's remote terminal.

Each policy enforcer 124, 126 preferably maintains a copy
of the self-extracting executable of the VPN client software
including a setup program and VPN reachability configuration
template. The setup program allows the VPN client software to
15 be installed on the VPN client 140. When downloading the self-
extracting executable, the configuration template is replaced
with the VPN reachability information that is specific to the
downloading user.

According to another embodiment of the invention, the system
20 allows the VPN client 140 to download a self-extracting
executable which, upon execution, only installs the VPN
reachability information that is unique to the user. According
to this embodiment, the VPN client software is already installed
on the VPN client 140. In this scenario, the setup program
25 allows the installation of the reachability information that is
specific to the downloading user, on the VPN client 140.

According to a third embodiment of the invention, the system
allows the VPN client 140 to automatically download the VPN
reachability information each time it connects to the policy
30 enforcer 124, 126. Thus, VPN reachability information is kept
up-to-date for each VPN client 140. Once a VPN session is
established, the connection between the VPN client 140 and the
policy enforcer is assumed to already be secure. The VPN client
preferably makes a common gateway interface (CGI) query to a web
35 server running on the policy enforcer, and downloads the current

5 VPN reachability information from the corresponding LDAP
database.

FIG. 16 is a block diagram of components in a self-extracting executable 290 according to one embodiment of the invention. The self-extracting executable 290 may be created
10 using commercially available tools such as the INSTALLSHIELD
EXEBUILDER of InstallShield Software Corporation of Schaumburg,
Illinois.

15 The self-extracting executable 290 preferably includes an
executable setup file 292 for installing the VPN client software
and/or the VPN configuration information. The setup file 292
preferably forms a static portion 298 of the self-extracting
executable since this information does not change based on the
downloading VPN client. The self-extracting executable 290
further includes VPN configuration file templates for the VPN
20 reachability information 294 and the VPN client's preshared key
information 296. The VPN reachability information 294 and the
VPN client's preshared key 296 preferably form a dynamic portion
299 of the self-extracting executable 290 since this information
changes based on the downloading VPN client. The self-extracting
25 executable 290 is then saved as a template file in the policy
enforcers 124, 126 and is ready to be downloaded by the remote
users.

FIG. 17 is a functional block diagram for downloading the
self-extracting executable 290 of FIG. 16 according to one
30 embodiment of the invention. In step 320, a new VPN client 140
first establishes a secure communication session with the policy
enforcer 124, 126 to download the self-extracting executable 290.
Preferably, this is accomplished via an HTTPS protocol session
on the VPN client's web browser or the like. In steps 322 and
35 324, the policy enforcer engages the VPN client in an
authentication procedure where the policy enforcer requests, and

5 the VPN client provides, his or her user name and password. In
step 326, the policy enforcer compares the provided information
with entries in its VPN client database 328. If the information
is correct, the policy enforcer finds appropriate preshared keys
for the user, and in step 330, also determines the VPN
10 reachability information of the client from a VPN configuration
database 332. The VPN client database 328 and VPN configuration
database 332 may reside as part of a single LDAP database 312,
314 managed by the policy enforcer 124, 126, or may constitute
separate LDAP databases.

15 In step 334, the policy enforcer replaces the dynamic
portion 299 of the self-extracting executable 290 with the VPN
reachability information and preshared key that is unique to the
VPN client. The newly generated self-extracting executable is
then downloaded to the VPN client 140 in step 336. When the
20 executable is run, it either installs the VPN client software
and/or the VPN reachability information.

Similar techniques may also be used for downloading a new
and updated copy of the VPN configuration information to the VPN
client each time the client connects to the policy enforcer and
25 negotiates a session key. In addition, the user may obtain the
latest configuration of the VPN network by expressly requesting
the policy enforcer for such information. Thus, the VPN client
need not be reinstalled and reconfigured each time updates are
made to the VPN reachability information.

30

VI. INTEGRATED POLICY ENFORCER

According to one embodiment of the invention, the
functionalities of the policy enforcer 124, 126 for policy
35 enforcement are partitioned for effective hardware
implementation. However, it should be apparent to one skilled
in the art that some or all of the functionalities may be

5 implemented in software, hardware, or various combinations thereof.

FIG. 18 is a schematic block diagram of the policy enforcer 124, 126 illustrating the partitioning of the various functionalities according to one embodiment of the invention. 10 The policy enforcer includes an Internet protocol security (IPSec) engine 502 for performing security and authentication functions in implementing, for instance, virtual private networks. A stream table 506 assembles the packets passing through the policy enforcer into streams. A protocol 15 classification engine 508 decodes the protocols used in forwarding the packets. A policy engine 510 enforces policies for the packets based on the policy settings stored in the policy database 132, 134. A packet forwarding module 504 receives packets from the public Internet via the router 110 and buffers, 20 forwards, or drops the packets based on the policies being enforced. A bandwidth management module 514 provides bandwidth shaping services to the packets being forwarded based on the bandwidth settings stored in the policy database 132, 134.

In practice, an incoming packet is matched against the 25 stream table 506 for determining if a matching entry already exists in the table. If not, a new entry is added. The stream table preferably includes enough portions of the packet to uniquely identify a stream. For example, in enforcing policies on IP layer three through layer four traffic, the stream table 30 may store a source IP, destination IP, source port, destination port, and protocol number of the incoming packet.

The protocol classification engine 508 takes the new stream and obtains a detailed protocol decode for the stream. The 35 policy engine 510 is then queried for the policy rules to be applied to the stream. Based on the policy rules returned by the policy engine 510, the packet forwarding module 504, IPSec engine

5 502, and/or the bandwidth management module 514 process the streams accordingly. The processing may be recursive until the packets in the stream have had all the actions specified by the policy rule set applied to them.

10 The policy enforcer also includes a statistics module 512 for collecting statistics on the packets forwarded through the local network as well as other status and resource usage information, and provides the same in logs and archives for sending to the policy server 122. According to one embodiment of the invention, the statistics module 512 keeps running byte 15 counts of the packets passing through the network 102, 104. These byte counts may be automatically sorted by classes, such as classes based on certain resources (e.g. users, hosts, services), as well as by bytes that are blocked by policies and exceptions, such as firewall policies. In this regard, the 20 statistics module 512 maintains in a cache a state table including a list of resources involved for each connection allowed through the firewall. For every packet flowing through the connection, the statistics module increments the packet and byte count for each of the resources in the list. The statistics 25 module 512 then forwards the organized information to the policy server 122 which enters the information directly into tables organized by classes and aged out periodically.

30 FIG. 19 is a more detailed schematic block diagram of the policy engine 510 according to one embodiment of the invention. The policy engine 510 includes a policy request table 602 that acts as a queue for all the policy decision requests. In this 35 regard, the portion of the packet matching the information stored in the stream table 506 is presented to the policy engine 510 in the form of a policy request. The policy request is then queued in the policy request table 602.

5 A resource engine 604 maintains an up-to-date mapping of
resource group names to member mappings. A policy rules database
buffer 608 stores a current policy rule set to be applied by the
policy engine 510. The policy rules stored in the buffer 608 are
preferably in the original group-based rule specification format.
10 Thus, the buffer 608 stores a rule created for a group in its
group-based form instead of instantiating a rule for each member
of the group.

15 A decision engine 606 includes logic to serve the incoming
policy decision requests in the policy request table 602 by
matching it against the policy rule set in the policy rules
database buffer 608 based on the actual membership information
obtained from the resource engine 604. The relevant group-based
rule matching the traffic is then identified and decision bits
in the stream table set for enforcing the corresponding actions.
20 The decision bits thus constitute the set of actions to be
performed on the packets of the stream. All packets matching the
streams are then processed based on these decision bits. The
decision engine may also specify an access control list (ACL)
including a set of rules that allow/deny traffic, a DiffServ
25 standard for providing a quality of service level to the traffic,
and/or VPN implementation information.

30 FIG. 20 is a more detailed schematic block diagram of the
protocol classification engine 508 according to one embodiment
of the invention. As illustrated in FIG. 20, the protocol
classification engine 508 includes a stream data assembly 702,
a sliding stream data window 704, an ASN.1 block 706, a protocol
classification state machine 708, and a protocol definition
signature database 710. The stream data assembly 702 extracts
35 and re-assembles the data portion of an input packet stream and
stores it in the sliding stream data window 704. Preferably, the
sliding stream data window follows first-in-first-out protocols.

5 The ASN.1 decoder further decodes the data stream, if needed, per conventional ASN.1 encoding/decoding standards. The protocol classification state machine 708 then matches the fully re-assembled and decoded data against the protocol definition 10 signature database 710. This database 710 preferably holds a mapping of protocol names to data patterns to be found in the data stream. The matched protocol is then returned to the stream table 506.

Thus, the protocol classification engine 508 provides extensive layer three through layer seven protocol decode and 15 packet classification, including complete identification of dynamic streams using a dynamically updated signature database compiled from scripted protocol definitions. As new protocols are defined in the future and/or users create their own custom 20 applications with custom protocols, a need may arise to add recognition of these protocols to the protocol classification engine. The described protocol classification engine architecture allows such additions by simply adding a new 25 scripted definition of the new protocol to the protocol classification engine without having to change the design each time a new protocol is added. This allows for custom protocol support and future protocol extensibility.

FIG. 21 is a more detailed schematic block diagram of the 30 IPSec engine 502 according to one embodiment of the invention. As illustrated in FIG. 21, the IPSec engine 502 includes a Pseudo-Random Number Generator (PRNG) function 802 for generating random numbers used for cryptographic key generation according to well known methods. A Diffie Hellman 804 and RSA 812 blocks implement the corresponding asymmetric public key 35 encryption/decryption/signature algorithms which are also well known in the art. An IKE block 806 communicates with an IPSec SA table 808 for implementing standard ISAKMP/Oakley(IKE) key

5 exchange protocols. A cryptographic transforms block 814 implements standard symmetric encryption/decryption algorithms. An IPSec Encapsulation/Decapsulation block 810 performs standard encapsulation/decapsulation functions. Accordingly, the IPSec engine 502 provides mature, standards-based IKE/IPSec
10 implementation with public key certificate support and necessary encryption/decryption functionality for packets passing through the private local networks 102, 104.

VII. NETWORK POLICY LOGS AND STATISTICS AGGREGATION

15 Referring again to FIG. 3, the log collecting and archiving module 304 collects information about the status and usage of resources from the policy enforcers 124, 126 as well as from the management module 302, and stores them in the archive database
20 318. The policy server reports module 316 then uses the collected logs and archives to generate reports in an organized report format.

25 According to one embodiment of the invention, each policy enforcer 124, 126 maintains a log file with information collected about the flow of traffic through the policy enforcer as well as the status and usage of resources associated with the policy enforcer. All the log files follow a predefined common log format, preferably designed to create compact logs.

30 FIG. 22 is a schematic layout diagram of such a log format according to one embodiment of the invention. Each log entry includes a timestamp 820 in the format yyyyymmddhhmmss, indicative of the year, month, date, hours, minutes, and seconds in which the log entry was created. A service field 822 indicates the type of service rendered by the policy enforcer 124, 126. Such
35 services include VPN, FTP, Telnet, HTTP, packet filter, bandwidth, and the like. Each log entry further includes a source IP address and port 824 indicating the source from where

5 a packet was received, as well as a destination IP address and port 826 indicating the destination to which the packet was forwarded.

10 A user ID field 828 identifies the user transmitting the packet. The user ID may be mapped to an entry in the LDAP database 130, 132, or 134 for obtaining additional details about the user.

15 A status field 830 indicates the status of an operation and may include a result code, error code, and the like. For example, for a packet filter service, the status field may include a result code "p" if the packet was passed or code "b" if the packet was blocked.

20 An operation field 832 indicates codes for a type of operation conducted by the service. For instance, operations for a VPN service may include sending packets and receiving packets. Operations for an FTP service may include GET and PUT operations. Operations for an HTTP service may include GET and POST operations.

25 In addition to the above, each log entry includes an in-bytes field 832 indicative of the number of bytes the policy enforcer received as a result of the activity, and an out-bytes field 834 indicative of the number of bytes transferred from the policy enforcer. Furthermore, a duration field 836 indicates the duration (e.g. in seconds) of the activity.

30 Certain fields of a particular log entry may be left blank if not applicable to a particular service. For instance, for an FTP download. Where there is no outgoing traffic, the out-bytes field is left blank. Furthermore, additional fields may be added based on the type of service being logged. For instance, for an HTTP activity, the URL that is accessed is also logged in the log entry. The additional fields are preferably appended to the end 35 of the standard log format.

5 A person skilled in the art should recognize that additions, deletions, and other types of modifications may be made to the log format without departing from the spirit and the scope of the invention as long as the log format common to all the policy enforcers and is aimed in creating compact logs.

10 The log files created by the policy enforcers 124, 126 are transferred to the policy server 122 based on archive options set by the policy server. In this regard, the network administrator specifies a threshold size for the logs created by the policy enforcers upon selection of the policy server archive option 752 of FIG. 9. When the log file exceeds the specified size, it is sent to the policy server 122. Preferably, the logs are transferred to the policy server 122 at least once a day even if the threshold size has not been exceeded. The logs may also be archived locally at the policy enforcer if so specified by the network administrator.

15 Once the policy server 122 receives the logs, it is stored in the archive database 318 preferably taking the form of an SQL database. The policy server reports module 316 queries this database to generate reports for each policy enforcer 124, 126. 20 In addition, the logs may be exported in a format that may be interpreted by commercially available products such as WEBTRENDS, manufactured by WebTrends Corporation of Portland, Oregon.

25 The reports created by the reports module 316 include summary usage reports for the various resources including policy enforcers, users, services, hosts, and VPNs. For instance, the reports may include VPN summary reports, bandwidth summary reports, packet filter reports, and the like, for each policy enforcer.

30 The reports preferably show usage of each of the resources over a period time. The start and the end date for the report may be specified by the user. The user may further drill down

5 on the time dimension and on the resource dimension for viewing
specific times and specific resources. For instance, in creating
the packet filter reports, the user may indicate a start and end
time, source IP address, source port, destination IP address, and
destination port. All packets meeting these criteria are then
10 fetched from the archive database 318 and shown in a packet
report.

VIII. METHOD FOR SELECTIVE LDAP DATABASE SYNCHRONIZATION

15 According to one embodiment of the invention, the databases
130, 132, 134 in the unified policy management system of FIG. 1
are LDAP databases storing policy management information
including policies for firewall, VPNs, bandwidth, administration,
user records, network records, services, and the like. As
20 described above, the LDAP directory service model is based on
entries where each entry is a collection of attributes. Entries
are arranged in a tree structure that follows a geographical and
organizational distribution. Entries are named according to
their position in the hierarchy by a distinguished name (DN).

25 The policy server 122 preferably stores the policy
management information for all the policy enforcers in the policy
server database 130. This information is organized in the
databases 130 as one or more DNs with corresponding attributes.
Appropriate portions of the policy server database are then
30 copied to the policy enforcer databases 132, 134.

FIG. 23 is a block diagram of an LDAP tree structure
including an LDAP root 265 and a plurality of branches 264, 266,
268, 270. According to one example, the policy server 122
maintains in the policy server database 130 branches 264 and 266
35 with policy management information for all the policy enforcers
124, 126. Each of the policy enforcers 124, 126 also maintain
portions of the branches 264 and/or 266 in their respective

5 policy enforcer databases 132, 134 as sub-trees of the policy
server database 130. The portions of the branches maintained by
each policy enforcer 124, 126 preferably relates to the
configuration information for that policy enforcer as well as
some additional information about the other policy enforcers.
10 This additional information is used to communicate with the other
policy enforcers.

The policy server 122 may further maintain branch 268
storing information used only by the applications running on the
server and not shared with any of the policy enforcers 124, 126.
15 Likewise, policy enforcers 124, 126 may maintain a portion of
branch 268 containing information used only by the applications
on each of the policy enforcers and not shared elsewhere.
Typically, the data stored in branch 268 is dynamically generated
and used by the applications running on the corresponding server
20 or agent.

Branch 270 is preferably only included in the LDAP tree for
the policy server database 130 and stores logged policy
management changes that may be propagated to the policy enforcers
124, 126. Such changes may include, for example, addition,
25 deletion, or modifications of a user on a device, VPN cloud,
bandwidth policy, or firewall policy made by the network
administrator via the various graphical user interfaces described
above. Such changes result in the updating of the policy
database 130 where the corresponding DN of the LDAP tree is
30 added, deleted, or modified. The policy server 122 further
creates a log of the changes and stores them in branch 270 for
later distribution to the policy enforcers 124, 126.

FIG. 24 is a more detailed block diagram of branch 270 of
the LDAP tree of FIG. 23. The LDAP root 265 includes an ApplyLog
35 270a entry which in turn includes a user log entry 270b and a
device log entry 270c. The user log entries include specific

5 administrator log entries identified by specific DNs 270d for reflecting the changes made by the particular administrators. The device log entry 270c includes specific device log entries identified by specific DNs 270e reflecting the changes that are to be distributed to the particular policy enforces 124, 126.
10 Preferably, the changes made by the administrators are propagated to the policy enforcers 124, 126 upon actuation of an apply button such as the apply button 417 illustrated in FIG. 6.

15 FIG. 25 is a flow diagram for logging and propagating LDAP changes to the policy enforcers according to one embodiment of the invention. In step 420, a particular network administrator makes a policy setting change. According to one example, the administrator is administrator "adm" working in the domain "domain1," and the change is the addition of a new user on a device.

20 In step 422, the change made the administrator is reflected in the policy server database 130. In this regard, branches 264 and 266 of the LDAP tree are modified accordingly to reflect the change in the policy setting. Additionally, in step 424, the policy server 122 creates a log of the changes for the administrator for later processing and sending to the appropriate policy agent. In step 426, the policy server 122 updates the administrator's log DN 270d to reflect the change. In the above example and as illustrated in FIG. 24, if the log created is named "A_L1," the policy server 122 updates the DN 270d for "adm" at "domain1" to create an attribute "apply" 270f that has the value "A_L1" 270g. Other changes made by the administrator are reflected in separate logs (e.g. "A_L2," "A_L3") and appended to the existing value of the apply attribute in the administrator's log DN 270d.

35 In step 428, the policy server 122 checks whether the changes made by the administrator are to be propagated to the

5 appropriate policy enforcers 124, 126. As discussed above, the changes are preferably propagated upon actuation of an apply button from the administrator's graphical user interface.

10 If the apply button has been actuated, the policy server creates, in step 430, a log for each policy enforcer to whom the change is to be transmitted. In this regard, the policy server 122 collects all the changes made by the administrator as reflected in the values 270g, 270h of the apply attribute 270f of the administrator's log DN 270d. These changes are processed for each policy enforcer belonging to the administrator's domain.

15 Such processing preferably involves picking the relevant changes and suitably modifying the DNs for the policy enforcer's LDAP. Such suitable modifications may be necessary, for instance, due to the differences in the tree structures in the policy server database 130 and the policy enforcer databases 132, 134. For instance, a change in the administrator's log may contain a DN that specifies the domain name of the policy enforcer. In applying this change to the policy enforcer, the domain name would not be specified in the DN since the policy enforcer's tree structure does not include a domain name.

20 25 The changes suitably modified for each policy enforcer's LDAP are then stored in a device log. Each policy enforcer's log DN 270e is then modified to reflect the change to the transmitted to the particular policy enforcer. In the above example and as illustrated in FIG. 24, if the device log created is named "PE_L1," the policy server 122 updates the DN 270e for the particular policy enforcer "PE1" at "domain1" to create an attribute "apply" 270i that has the value "PE_L1" 270j.

30 35 In step 432, the apply attribute 270f for the administrator's log DN 270d is then deleted from the LDAP tree. In step 434, the changes collected for each policy enforcer, as reflected in the values 270j, 270k of the apply attribute 270i

5 of the policy enforcer's log DN 270e, are transmitted to the policy enforcer for updating its database 132, 134. The changes are sent to the policy enforcers preferably over the HTTPS channel.

10 In step 436, the policy server 122 checks whether the updates have been successful. In this regard, the policy server 122 waits to receive an acknowledgment from the policy enforcer that the updates have been successfully completed. Upon a positive response from the policy enforcer, the policy server 122 deletes the apply attribute 270e for the policy enforcer's log 15 DN 270e. Otherwise, if the update was not successful (e.g. because the policy enforcer was down), the apply log is re-sent the next time another apply function is invoked. Alternatively, the failed policy enforcer transmits a request to the policy server 122 of the log of non-applied changes when it rejoins the 20 network (e.g. by rebooting).

IX. STATE TRANSITION PROTOCOL FOR HIGH AVAILABILITY UNITS

25 According to one embodiment of the invention, the policy server 122, policy enforcers 124, 126, as well as other network devices may be configured for high availability by maintaining a backup unit in addition to a primary unit.

30 FIG. 26 is a schematic block diagram of a high availability system including a primary unit 902 and a backup unit 904. The two units 902, 904 communicate with each other by exchanging heartbeats over parallel ports 906a, 906b and a cable 908. Such parallel ports 906a, 906b and cable 908 are conventional components that are commonly available in the art.

35 The primary unit 902 and the backup unit 904 are each similarly connected to other components 910, 912, 914 via ports 920a, 920b, 922a, 922b, 924a, 924b, respectively. These components 910, 912, 914 may be hubs, switches, connectors, or

5 the like. Because the primary unit 902 and backup unit 904 provide similar services and functions and may be used interchangeably, each unit is preferably connected to the same components 910, 912, 914.

10 The parallel port cable 908 is preferably a conventional laplink cable designed to connect two parallel ports and allow communications between them. The primary unit 902 and the backup unit 904 preferably communicate with each other via TCP packets over the high-availability ports 906a, 906b. A point-to-point connection preferably exists between the primary unit 902 and the 15 backup unit 904 over the high-availability ports 906a, 906b.

20 The primary unit 902 is preferably responsible for checking the status of its network ports for problems or failures. For example, if the primary unit 902 detects that one of its network ports is inoperable, e.g. port 922a, the primary unit 902 then checks whether the corresponding port 922b in the backup unit 904 is operational. Upon determining that the corresponding port 922b in the backup unit 904 is operational, the primary unit 902 sends a request to the backup unit 904 to take over the system 25 functions as the active unit. The primary unit 902 then relinquishes its role as the active unit and shuts itself down, allowing the backup unit 904 to take on the responsibilities of the primary unit 902. When the primary unit 902 restarts operation, the backup unit 904 receives a request from the primary unit 902 to relinquish its role as the active unit.

30 When the primary unit 902 is active and does not detect any defects in its ports, it continuously listens on the high-availability port 906a to keep track of the status of the backup unit 904. The primary unit 902 continues to listen on the high-availability port 906a for signals coming from the backup unit 35 904. When the backup unit 904 is up and running, it connects to the primary unit 902. Once the connection is made, the backup

5 unit 904 begins sending heartbeats to the primary unit 902. The
 backup unit 904 continuously sends heartbeats to the primary unit
 902 in predetermined intervals. According to one embodiment of
 the invention, the backup unit 904 sends a "Keep Alive" packet
 including a KEEP_ALIVE command to the primary unit 902 every one
10 second.

15 The primary unit 902 responds to the "Keep Alive" packet by
 changing the command field of the packet to a KEEP_ALIVE_RESP
 command and re-transmitting the packet to the sender. If the
 backup unit 904 does not receive a response back from the primary
15 unit 902 for a predetermined period of time (e.g. one second) for
 one "Keep Alive" packet, the backup unit 904 begins preparing to
 take over the active role. Preferably, the predetermined period
 should not be greater less than two consecutive "Keep Alive"
 packets.

20 Upon taking the role of the active unit, the backup unit 904
 attempts to reestablish a connection with the primary unit 902
 at regular intervals to determine whether the problem or failure
 in the primary unit has been cured. If the problem or failure
 has been cured, the backup unit 904 relinquishes its control to
25 the primary unit 902 after setting the IP addresses of all the
 network interface cards to the assigned value.

30 In situations where the backup unit 904 takes over the
 active role from the primary unit 902, an alert/alarm is sent to
 the network administrator indicating such a change. In addition,
 if the primary unit 902 does not receive heartbeats from the
 backup unit 904, an alert/alarm is sent to the administrator
 indicating that the backup unit has failed.

35 A situation may arise when both the primary unit 902 and the
 backup unit 904 are fully functional, and the backup unit 904
 desires to take over the active role. In this case, the backup
 unit 904 transmits a shut-down command to the primary unit 902

5 which then relinquishes control. The backup unit 904 continues its role as the active unit until the primary unit 902 transmits a request to the backup unit 904 to relinquish its active role.

According to one embodiment of the invention, the initial status determination protocol of each high availability unit as 10 a primary, backup, or stand-alone unit relies on a self-discovery process. FIG. 27 is a flow diagram of an exemplary status discovery process according to one embodiment of the invention. In step 930, a first high availability unit (unit X) that has not 15 yet definitively discovered its status as a primary or a backup unit boots up, and in step 932 assumes the role of a backup unit. In step 934, unit X searches the network for a primary unit and inquires, in step 936, whether a primary unit has been detected. If the answer is YES, unit X tries to connect to the primary 20 unit. If it is successful, unit X initializes as the backup unit in step 938. If, on the other hand, unit X does not detect the primary unit, unit X assumes the role of the primary unit in step 940.

In step 942, unit X searches the network for a backup unit. 25 If the backup unit is detected, as inquired in step 944, unit X connects to the backup unit and initializes as the primary unit in step 946. If, on the other hand, unit X does not detect any other units in the network within a predetermined time, unit X initializes as a stand-alone unit in step 948.

Once the primary and secondary units have been initialized, 30 configuration changes of the primary unit are also transferred to the backup unit in order to keep the two units synchronized. The configuration information is preferably stored in an LDAP database such as the central policy server database 130 or policy agent databases 124, 126.

35 FIG. 28 is a flow diagram of a process for maintaining configuration information synchronized in the primary and backup

5 units. In step 950, the primary unit boots up and in step 952, detects the backup unit. In step 954, the backup unit receives configuration change information from the primary unit if it is functional. Otherwise, the configuration changes are entered directly into the backup unit by the network administrator. If 10 the configuration change is to be received from the primary unit, the primary unit notifies the backup unit when configuration changes occur in the primary unit. The changes are then transferred and applied to the backup unit. The backup unit in turn transmits the status of the transfer and the apply back to 15 the primary unit.

In step 956, the primary unit is checked to determine whether it is functional. If it is, the primary unit is likewise updated with the configuration change. Otherwise, if the primary unit is not functional, the backup unit takes on the active role 20 and becomes the active unit in step 958. The primary unit may become non-functional and thus, inactive, due failures in the CPU board, the network interface card, or power supply.

In step 960, the backup unit tags the changes to transfer them to the primary once the primary becomes functional. Once 25 the primary unit becomes functional, the primary unit is updated with the tagged changes maintained by the backup unit as is reflected in step 962.

According to one embodiment of the invention, software updates on the primary and backup units are also synchronized so 30 as to update the primary and backup units serially in a single cycle without the need for multiple update cycles. Thus, the network administrator need not duplicate the efforts of updating the backup unit with the same information as the primary unit.

FIG. 29 is an exemplary flow diagram of updating the primary 35 and backup units when they are both functional. In step 970, an update, such as a software update not stored in the LDAP

5 databases, is sent/transmitted to the primary unit from a management station accessible by the network administrator. The primary unit then updates itself in step 972. In step 974, the primary unit automatically sends/transmits the update information to the backup unit. In step 976, the backup unit updates itself 10 with the update information received from the primary unit.

FIG. 30 is an exemplary flow diagram of updating the primary and backup units when the primary unit is not functional. In step 978, the primary unit becomes nonfunctional, and in step 980, the network administrator sends/transmits an upgrade 15 directly to the backup unit instead of the primary unit. In step 982, the backup unit updates itself with the information received from the management station and waits for the primary unit to become functional. Once the primary unit becomes functional, the update is automatically sent/transmitted to the primary unit for 20 upgrading in step 986. The primary unit then updates itself in step 988.

Although the present invention has been described in detail with reference to the preferred embodiments thereof, those skilled in the art will appreciate that various substitutions and 25 modifications can be made to the examples described herein while remaining within the spirit and scope of the invention as defined in the appended claims.

For example, the unified policy management system of FIG. 1 should be viewed as illustrative rather than limiting. It 30 should be apparent to those skilled in the art who are enlightened by the present invention that many alternative configurations are possible. For example, there may be additional networks with policy enforcers or no additional networks at all. Likewise, policy enforcers may not necessarily access the policy 35 server through the Internet, but may be connected via other means such as a WAN, MAN, etc. In short, the number and type of users

5 and resources within and without the organization can vary greatly while staying within the scope of the invention.

5 CLAIMS:

10 1. A system for managing policy services in an organization, the organization including a first network having a first set of resources and a second network remote from the first network having a second set of resources, the system comprising:

a first edge device associated with the first network, the first edge device configured to manage policies for the first network and the first set of resources in accordance with first policy settings stored in a first database;

15 a second edge device associated with the second network, the second edge device configured to manage policies for the second network and the second set of resources in accordance with second policy settings stored in a second database; and

20 25 a central policy server defining the first and second policy settings and managing the first and second edge devices from a single location, the central policy server being associated with a central database storing configuration information of the first and second edge devices, wherein the central database is organized according to a hierarchical object oriented structure.

30 2. The system of claim 1, wherein the first and second databases are organized according to the hierarchical object oriented structure.

3. The system of claim 1, wherein the configuration information includes the first and second policy settings.

35 4. The system of claim 3, wherein the hierarchical object oriented structure includes a plurality of resource objects and policy objects for defining the first and second policy settings.

5 5. The system of claim 4, wherein the central database and
the first and second databases are Lightweight Directory Access
Protocol (LDAP) databases storing each resource object and policy
object as an LDAP entry.

10 6. The system of claim 4, wherein the resource objects are
selected from a group consisting of devices, users, hosts,
services, and time.

15 7. The system of claim 6, wherein the devices include the
first and second edge devices, each device being associated with
a set of users and a particular host.

20 8. The system of claim 6, wherein the hosts include the
first and second networks.

25 9. The system of claim 4, wherein the policy objects are
selected from a group consisting of bandwidth, firewall,
administration, and virtual private network grouping.

30 10. The system of claim 9, wherein the virtual private
network grouping includes a virtual private network associated
with one or more sites, users, and rules.

35 11. The system of claim 10, wherein each site includes one
or more networks behind an edge device.

12. The system of claim 10, wherein the rules are firewall
rules providing access control over network traffic flowing
through the virtual private network.

5 13. In a system including a first network having a first
set of resources and a second network remote from the first
network having a second set of resources, the first network being
associated with a first edge device and a first database, and the
second network being associated with a second edge device and a
10 second database, the system further including a central policy
server in communication with the first and second edge devices,
the central policy server being associated with a central
database, a method for managing policy services in the system
comprising:

15 storing configuration information of the first and
second edge devices in the central database, the central database
being organized in a hierarchical object oriented structure;

 storing first policy settings in the first database;

 storing second policy settings in the second database;

20 managing policies for the first network and the first
set of resources from the first edge device in accordance with
the first policy settings stored in the first database;

25 managing policies for the second network and the second
set of resources from the second edge device in accordance with
the second policy settings stored in the second database; and

 defining the first and second policy settings and
managing the first and second edge devices from the central
policy server.

30 14. The method of claim 13, wherein the first and second
databases are organized according to the hierarchical object
oriented structure.

35 15. The method of claim 13, wherein the configuration
information includes the first and second policy settings.

5 16. The method of claim 15, wherein the hierarchical object oriented structure includes a plurality of resource objects and policy objects for defining the first and second policy settings.

10 17. The method of claim 16, wherein the central database and the first and second databases are Lightweight Directory Access Protocol (LDAP) databases storing each resource object and policy object as an LDAP entry.

15 18. The method of claim 16, wherein the resource objects are selected from a group consisting of devices, users, hosts, services, and time.

20 19. The method of claim 18, wherein the devices include the first and second edge devices, each device being associated with a set of users and a particular host.

25 20. The method of claim 18, wherein the hosts include the first and second networks.

25 21. The method of claim 16, wherein the policy objects are selected from a group consisting of bandwidth, firewall, administration, and virtual private network grouping.

30 22. The method of claim 21, wherein the virtual private network grouping includes a virtual private network associated with one or more sites, users, and rules.

23. The method of claim 22, wherein each site includes one or more networks behind an edge device.

5 24. The method of claim 22, wherein the rules are firewall
rules providing access control over network traffic flowing
through the virtual private network.

OBJECT MODEL FOR NETWORK POLICY MANAGEMENT**ABSTRACT OF THE DISCLOSURE**

A unified policy management system for an organization including a central policy server and remotely situated policy enforcers. A central database and policy enforcer databases storing policy settings are configured as LDAP databases adhering to a hierarchical object oriented structure. Such structure allows the policy settings to be defined in an intuitive and extensible fashion. Changes in the policy settings made at the central policy server are automatically transferred to the policy enforcers for updating their respective databases. Each policy enforcer collects and transmits health and status information in a predefined log format and transmits it to the policy server for efficient monitoring by the policy server. For further efficiencies, the policy enforcement functionalities of the policy enforcers are effectively partitioned so as to be readily implemented in hardware. The system also provides for dynamically routed VPNs where VPN membership lists are automatically created and shared with the member policy enforcers. Updates to such membership lists are also automatically transferred to remote VPN clients. The system further provides for fine grain access control of the traffic in the VPN by allowing definition of firewall rules within the VPN. In addition, policy server and policy enforcers may be configured for high availability by maintaining a backup unit in addition to a primary unit. The backup unit become active upon failure of the primary unit.

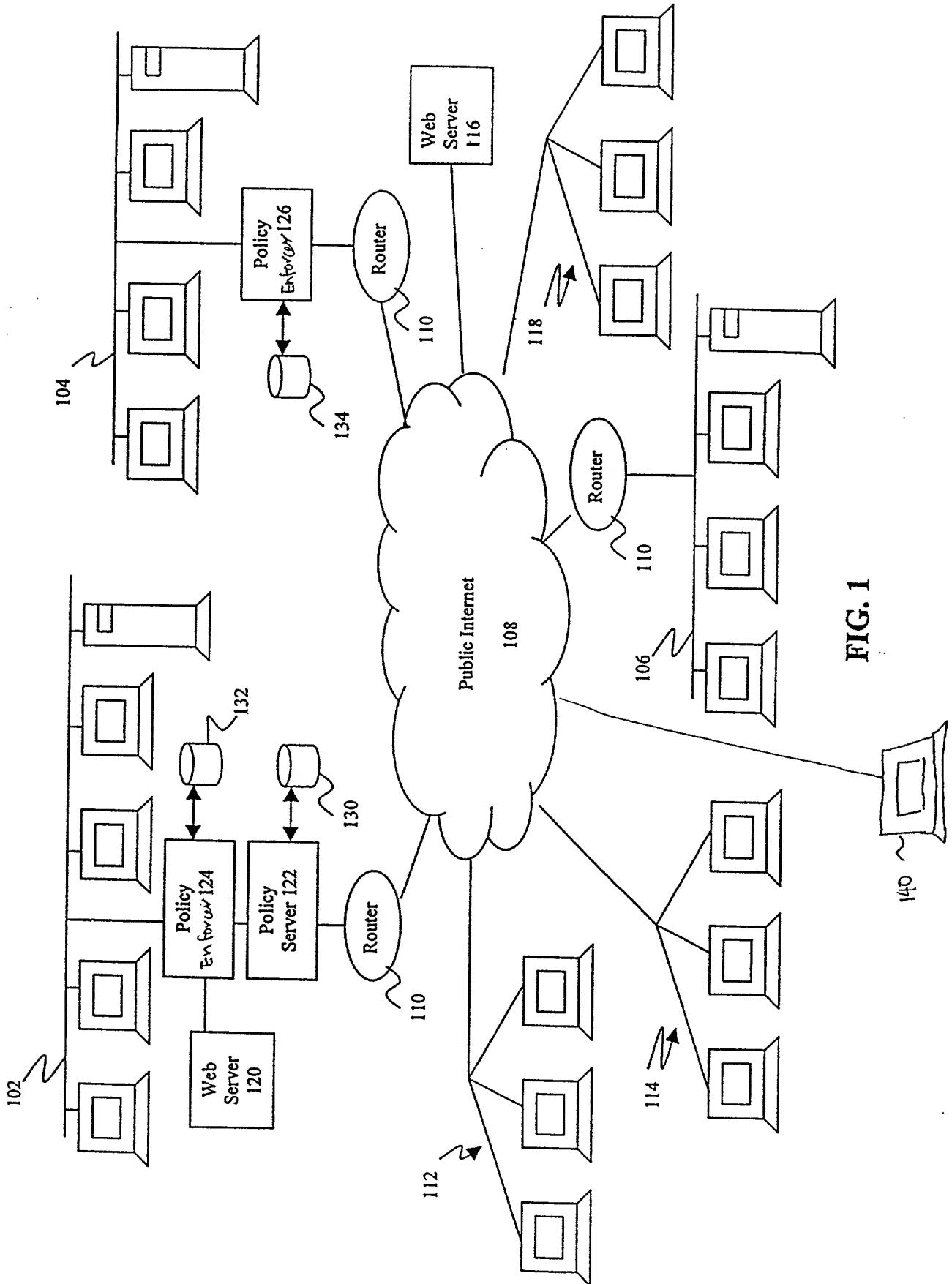
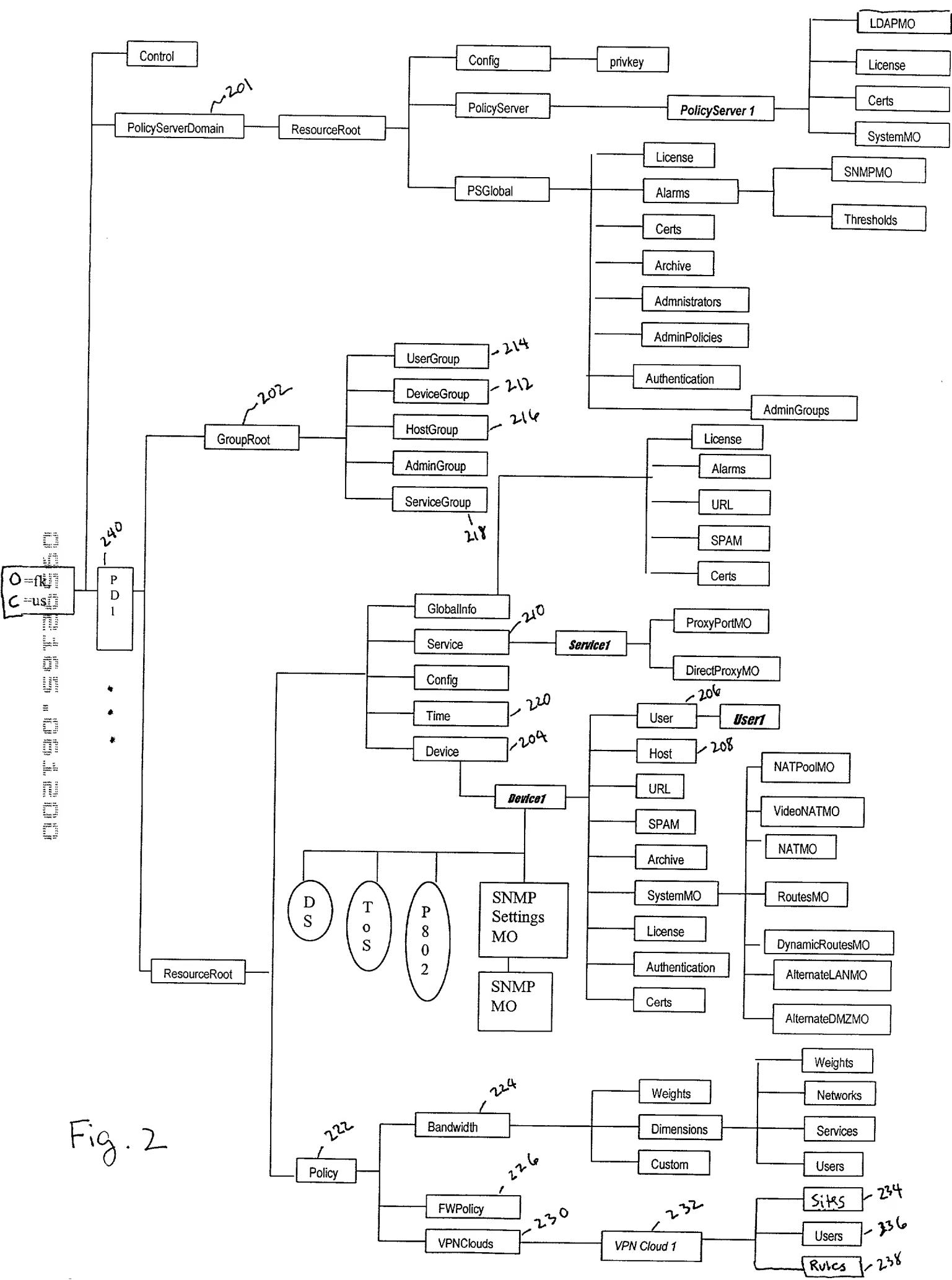


FIG. 1



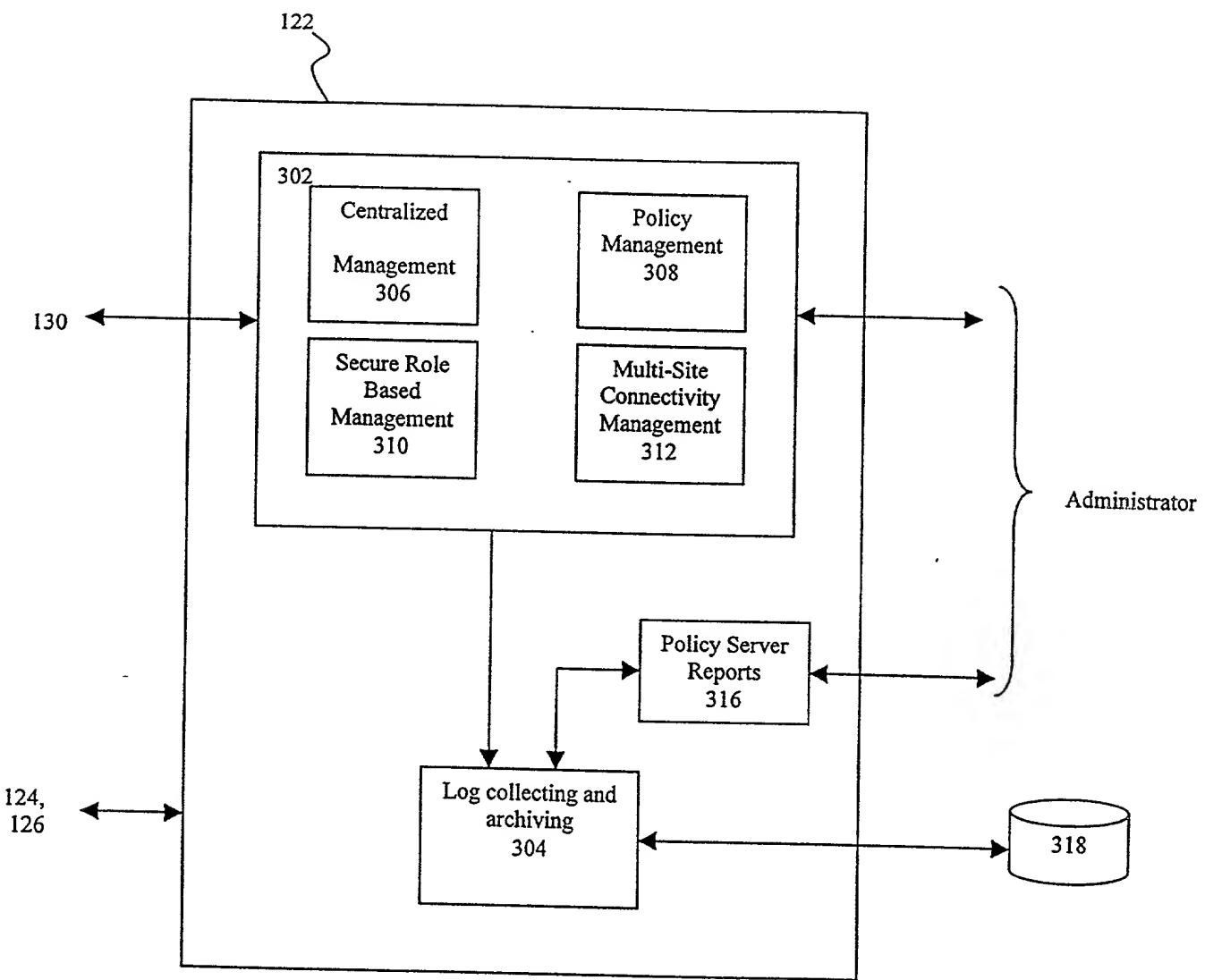


FIG. 3

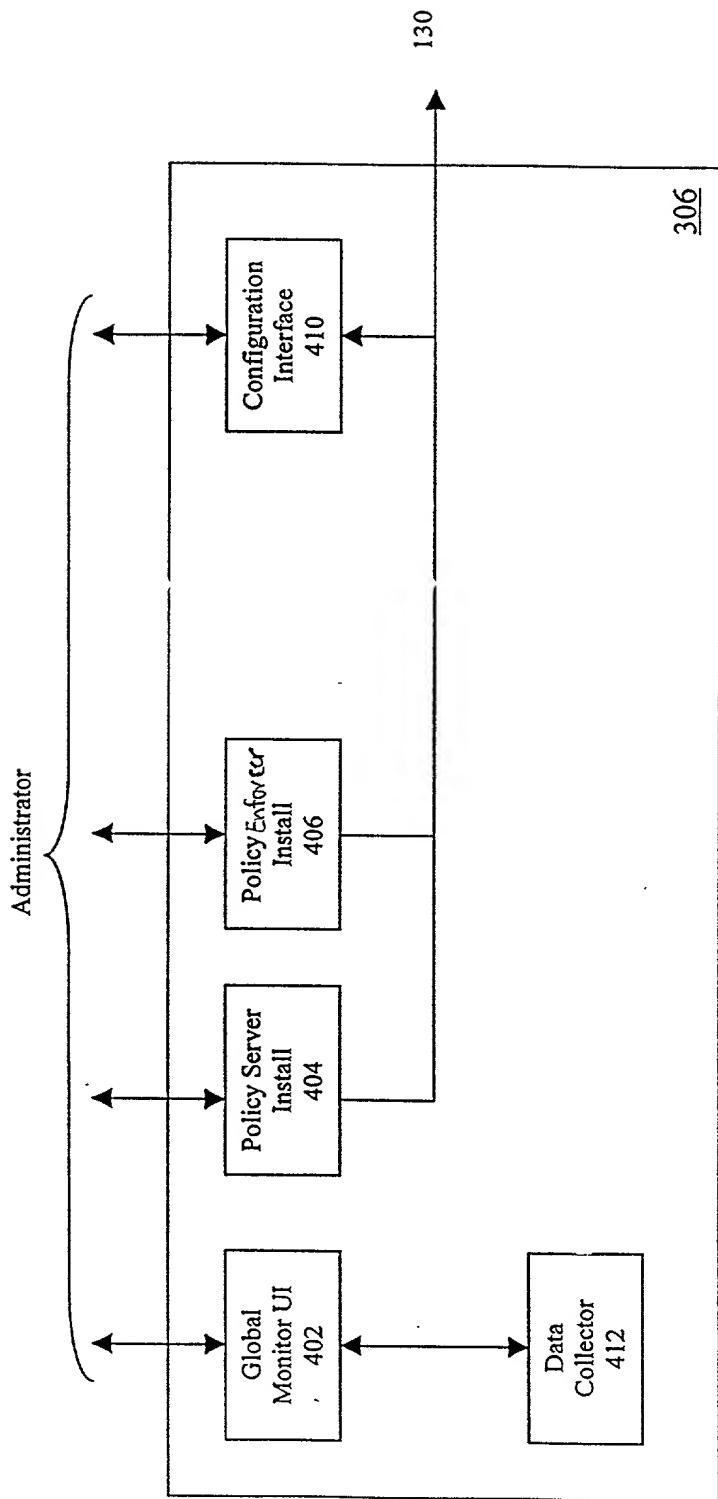


FIG. 4

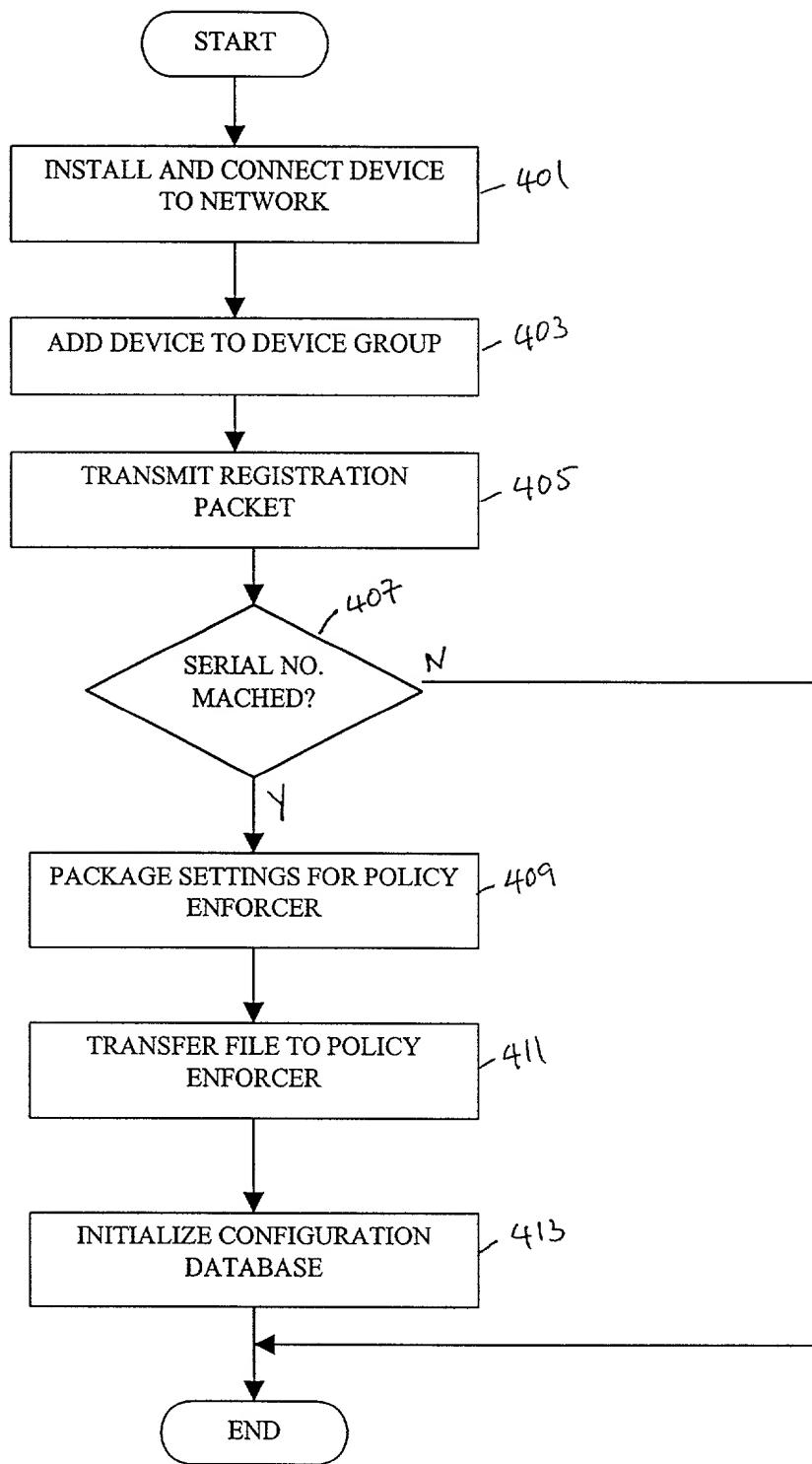


FIG. 5

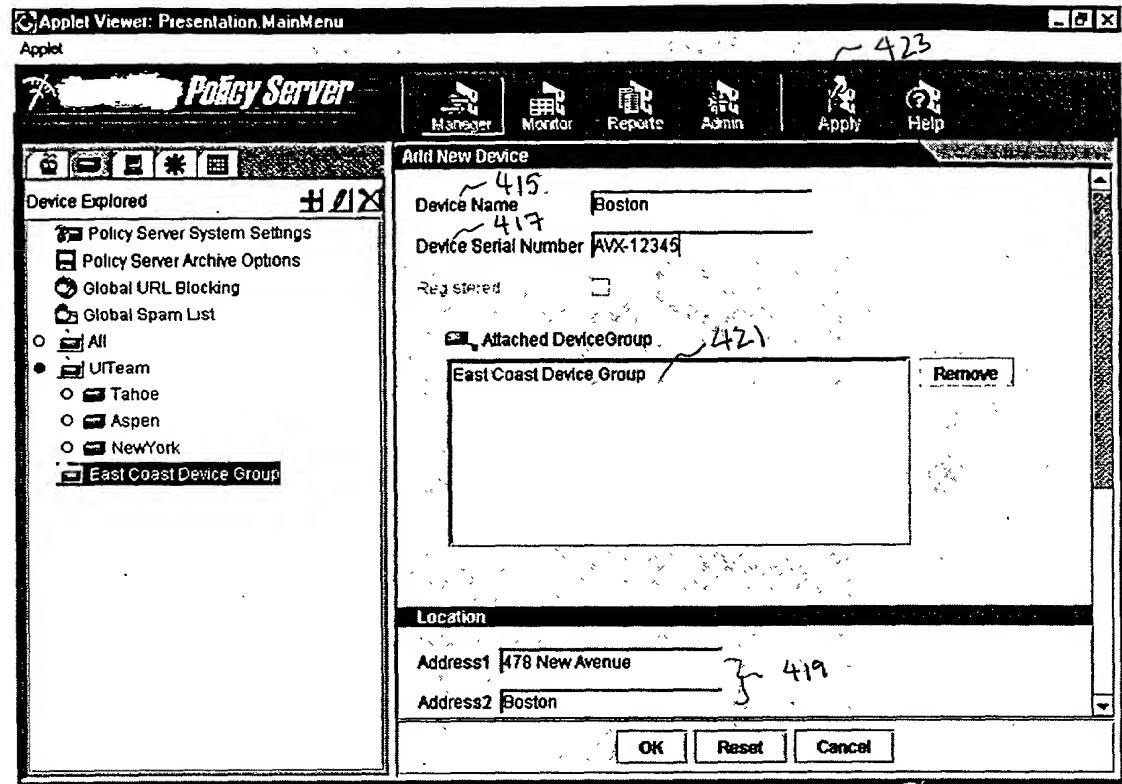


Fig. 6

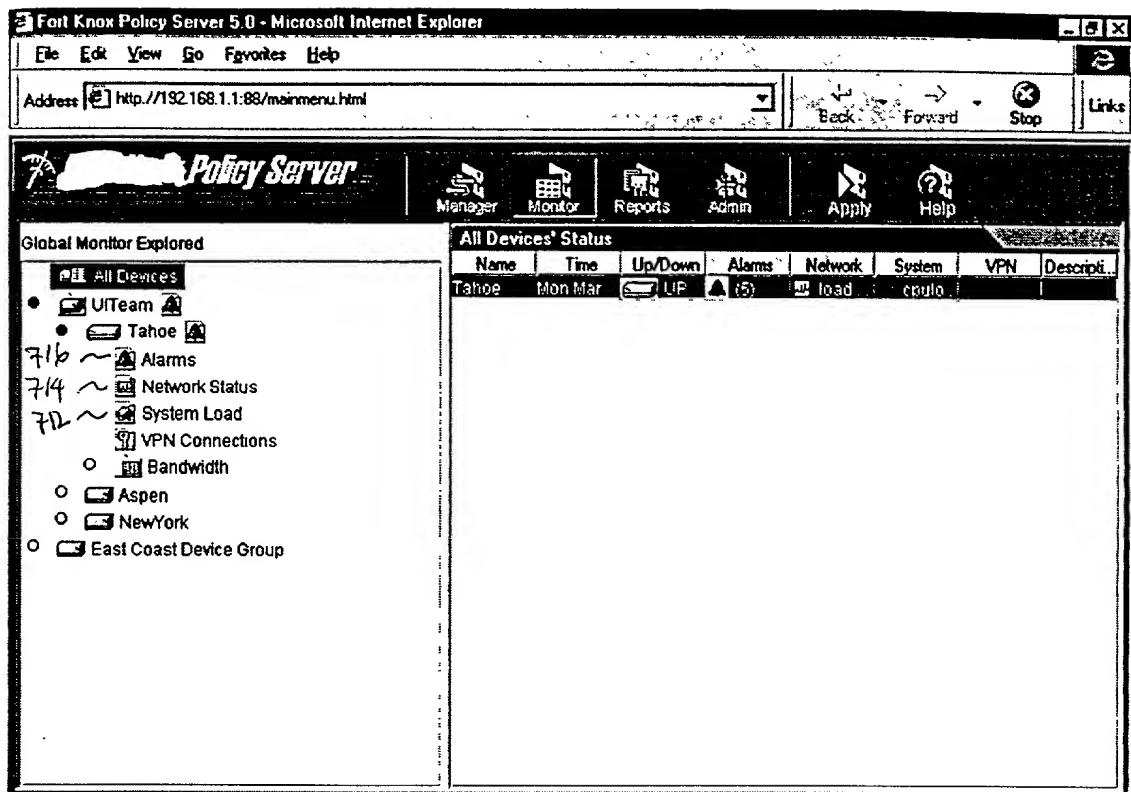


Fig. 7

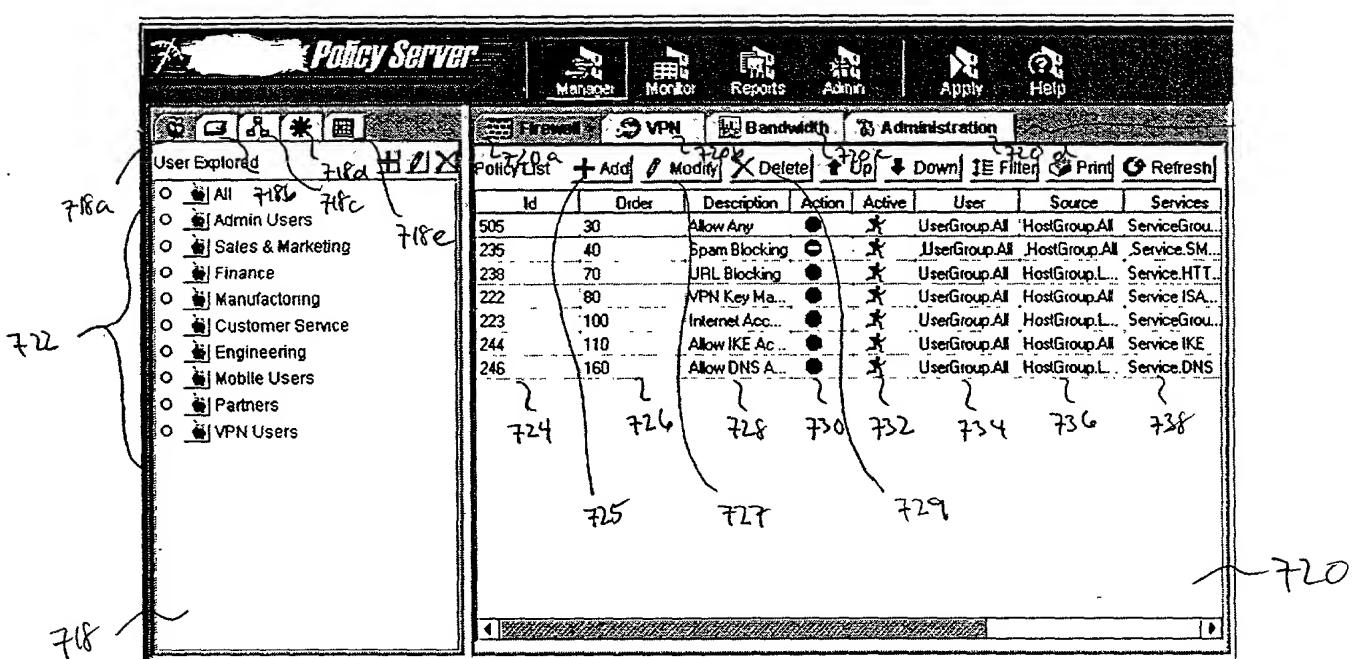


Fig. 8

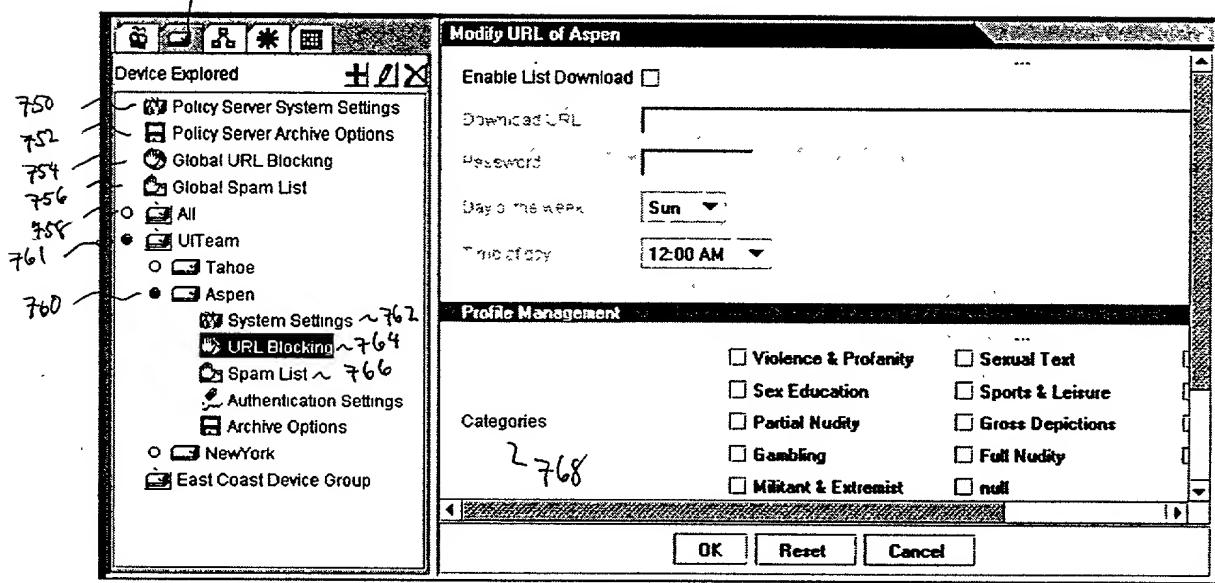


Fig. 9

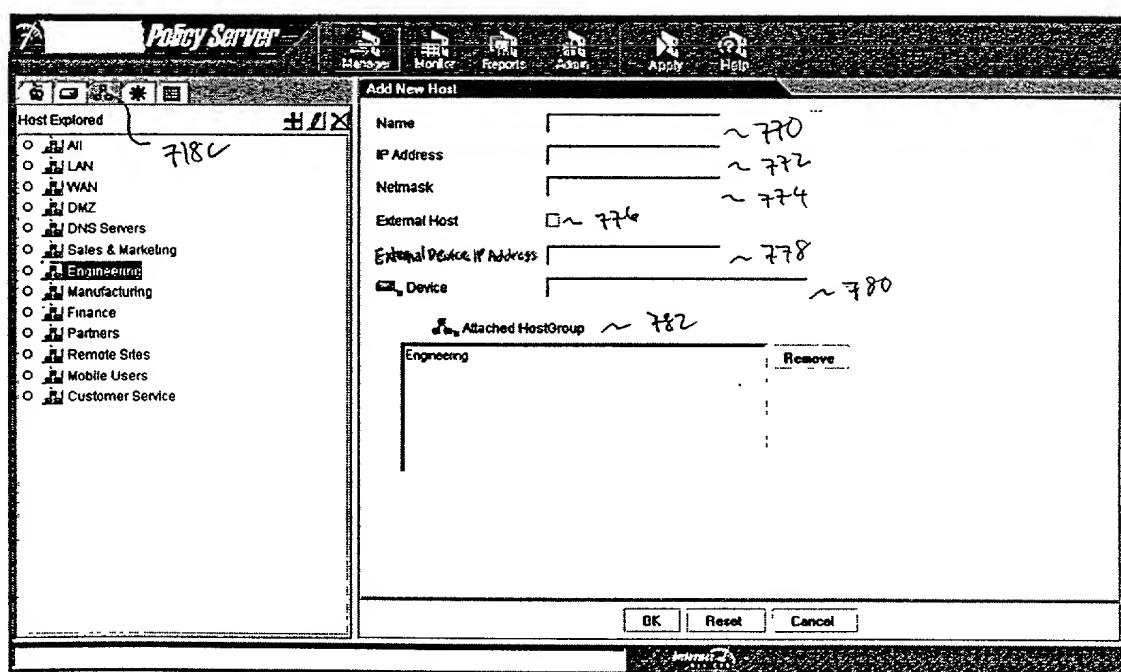


Fig. 10

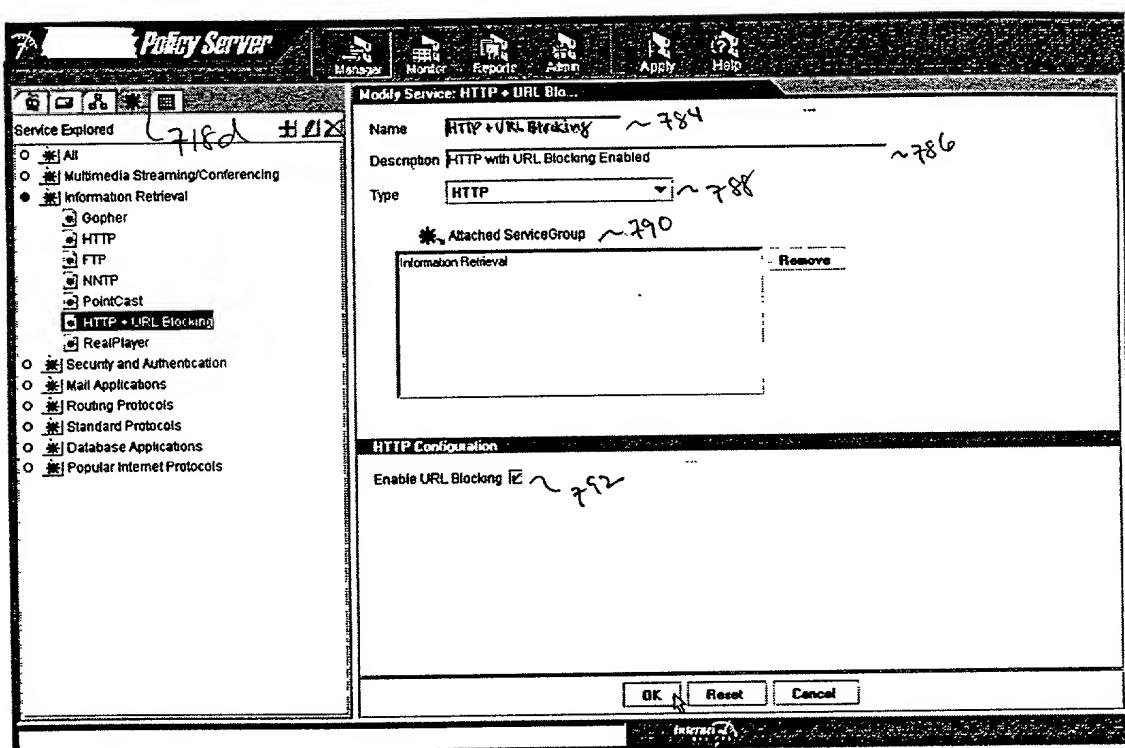


Fig. 11

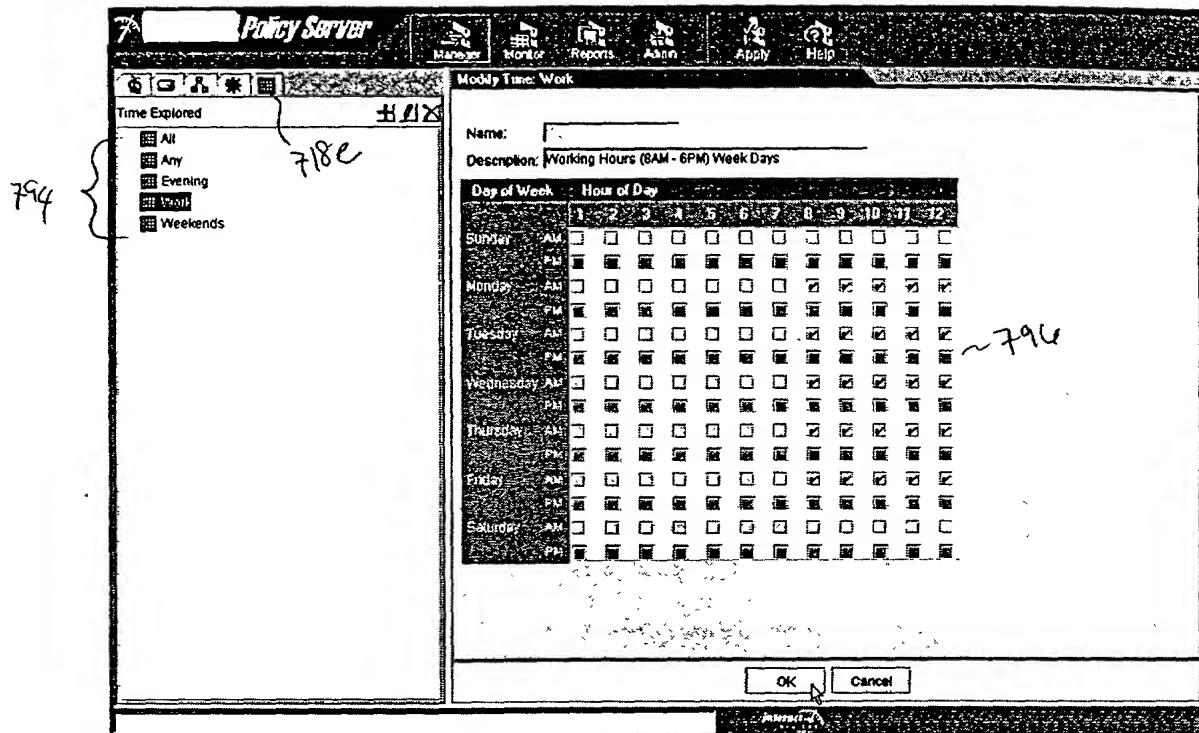


Fig. 12

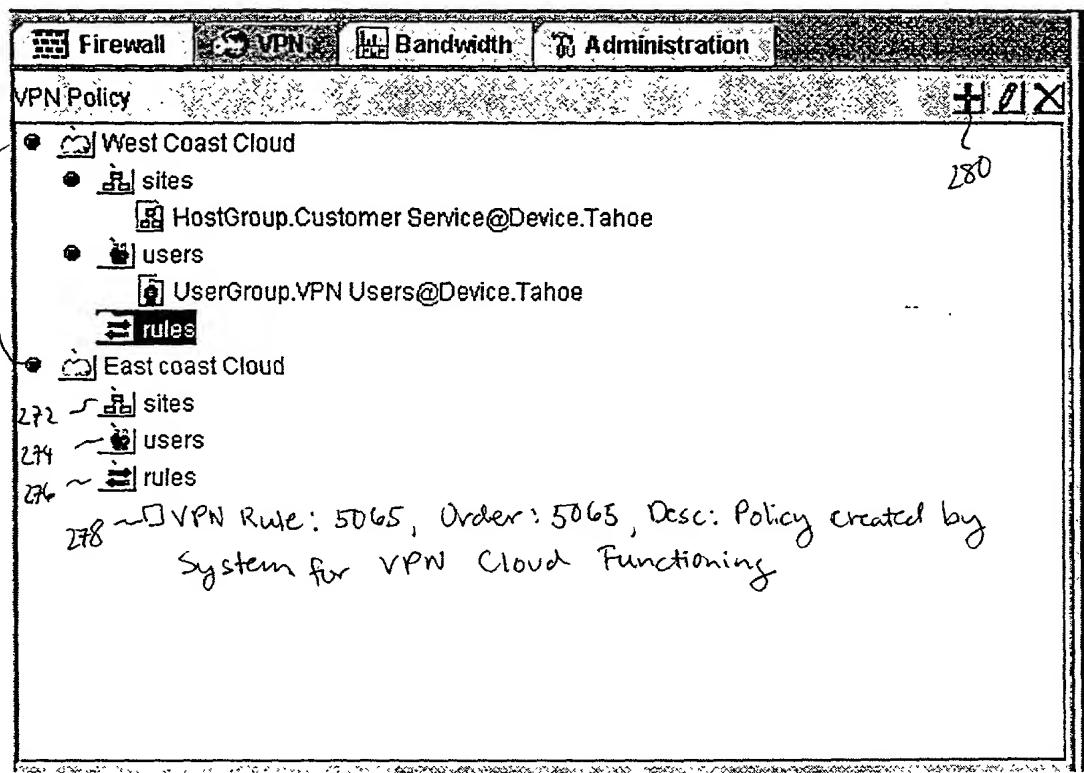


FIG. 13

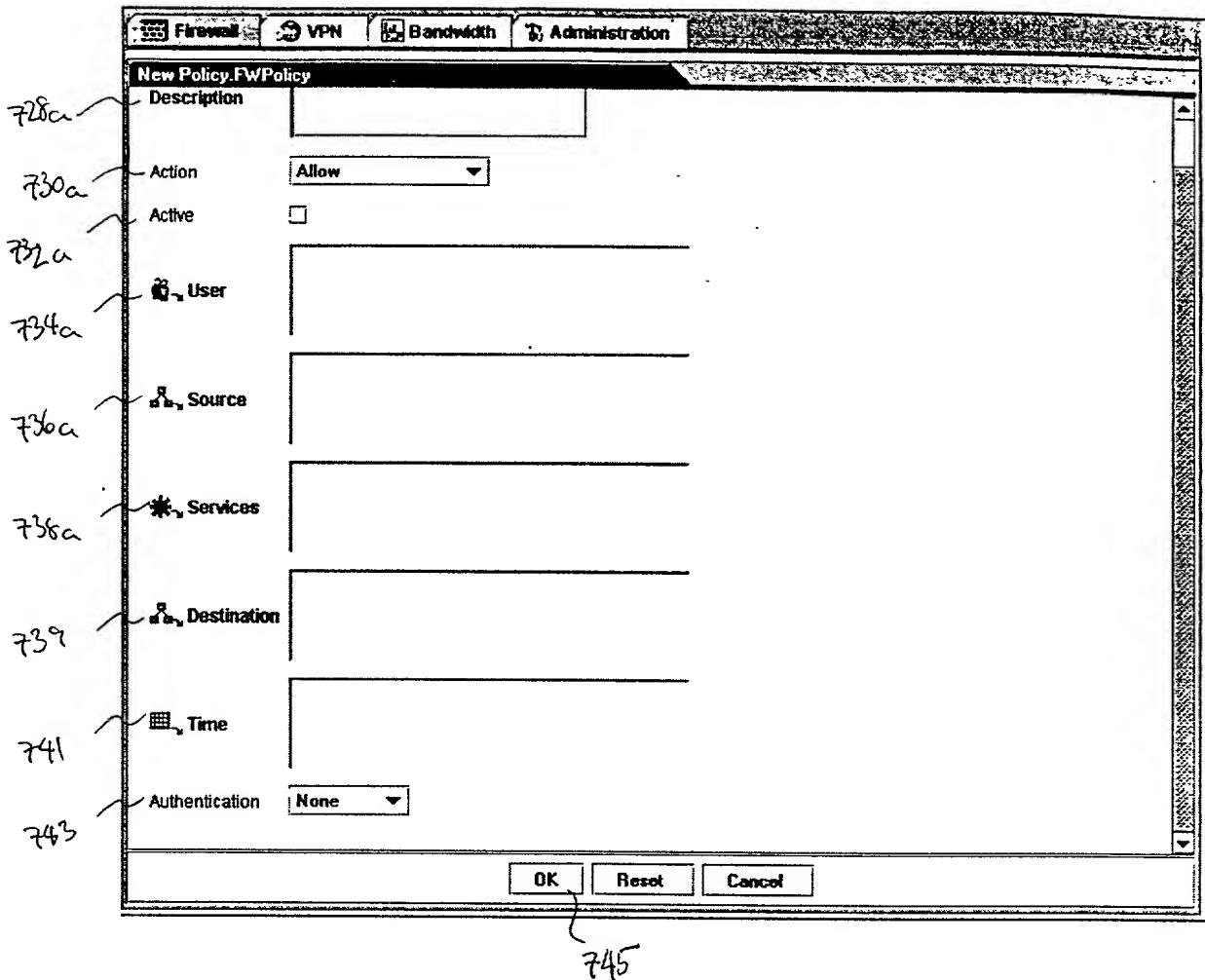


Fig. 14

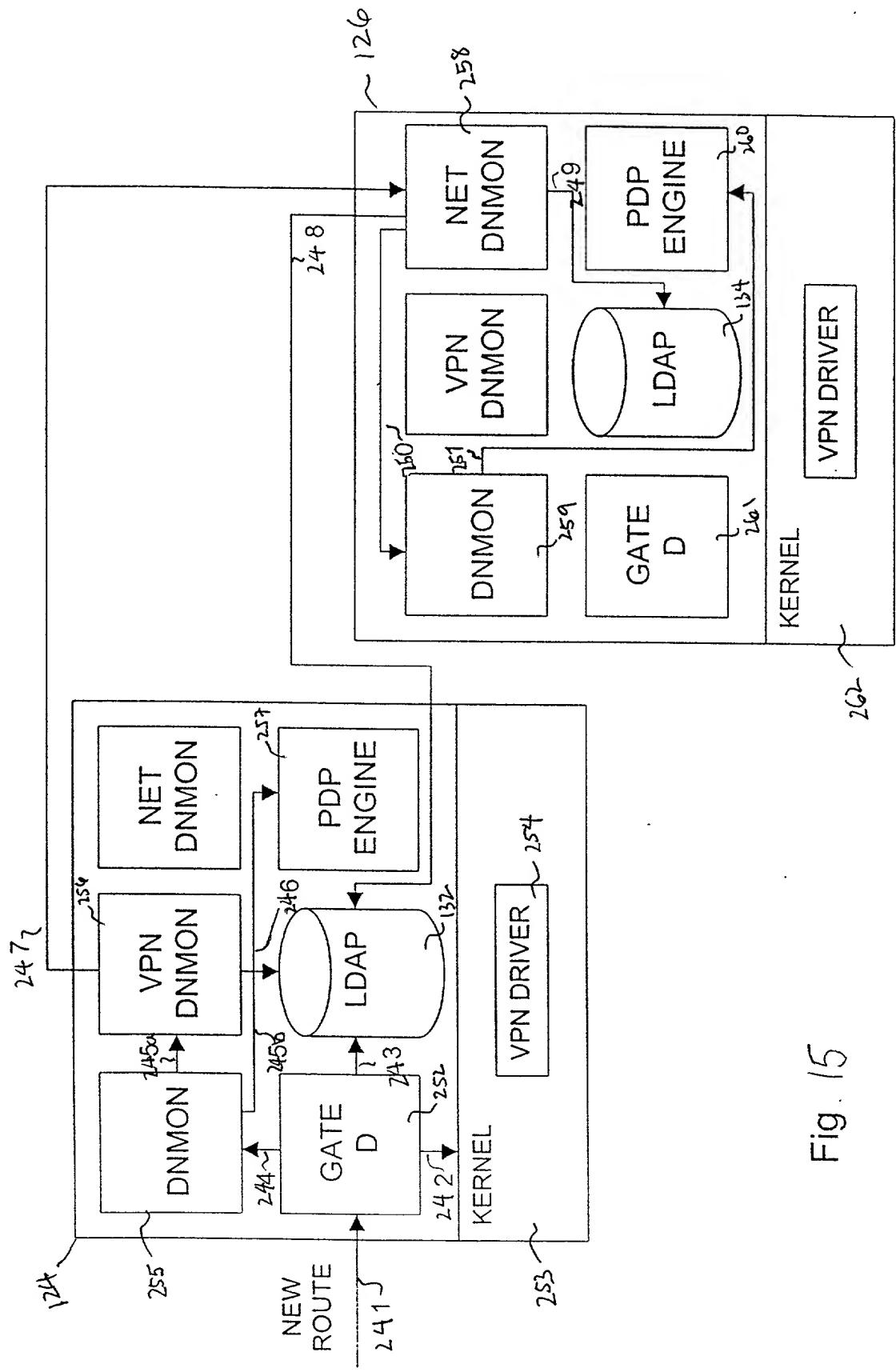


Fig. 15

FIG. 16

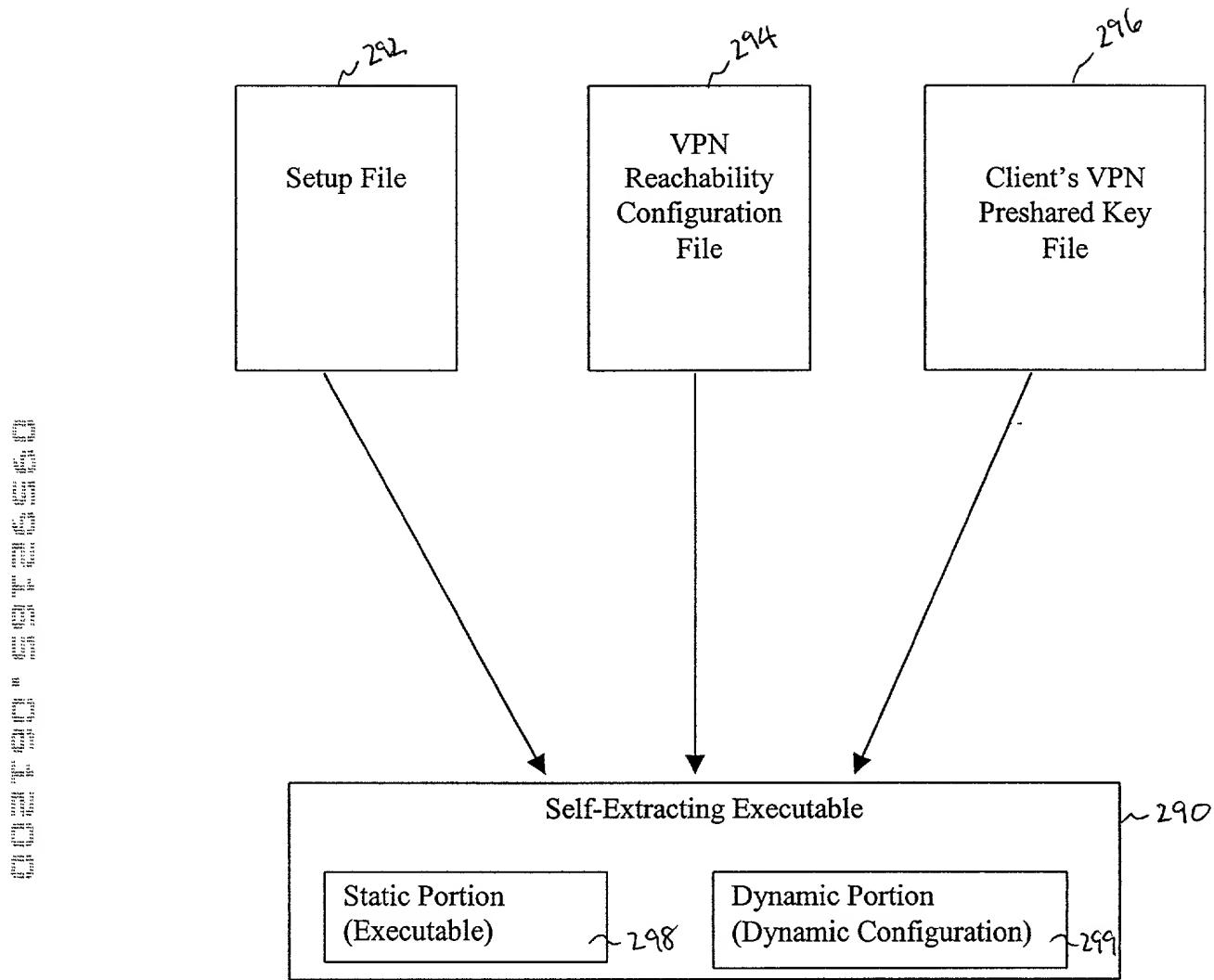
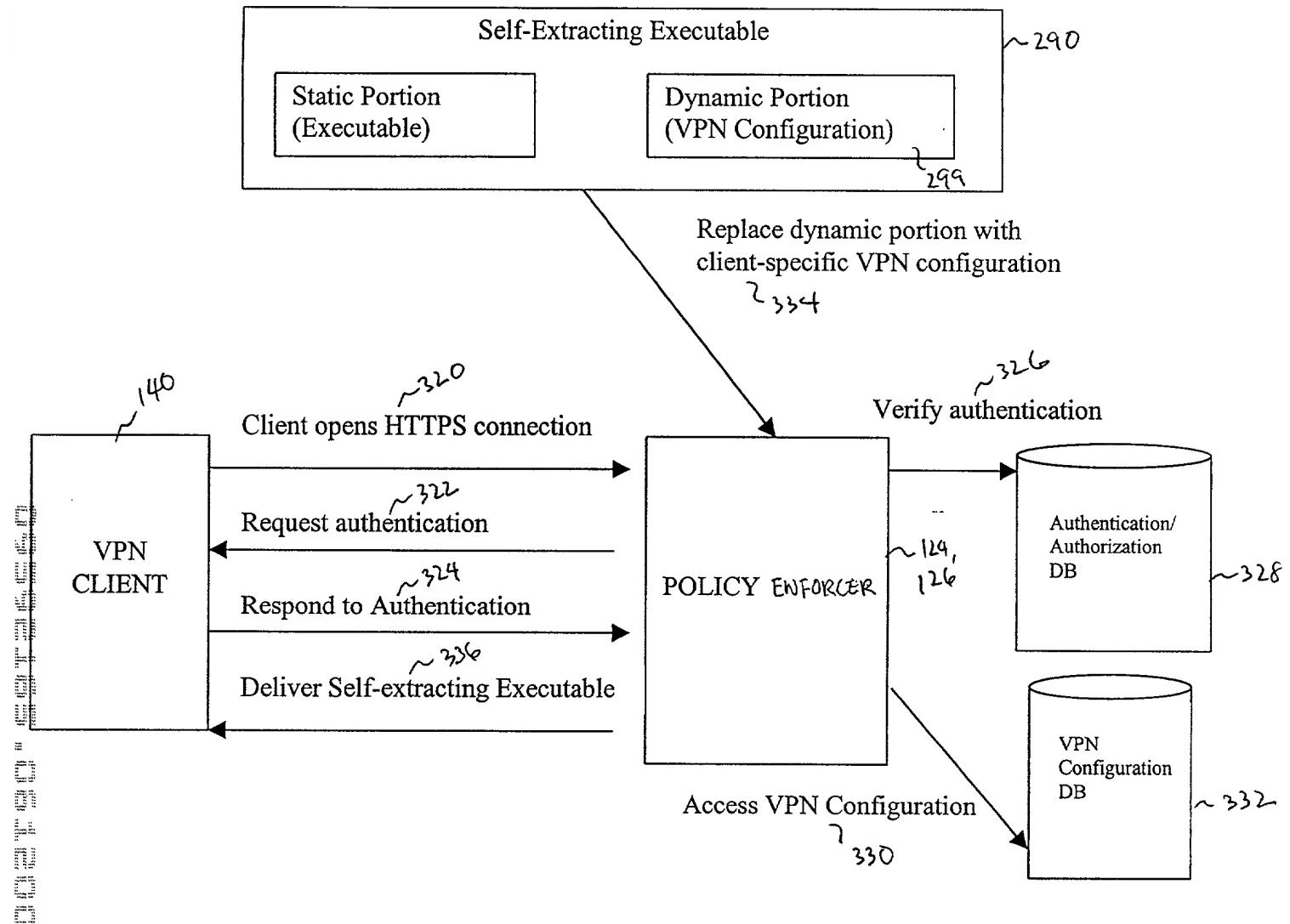


FIG. 17



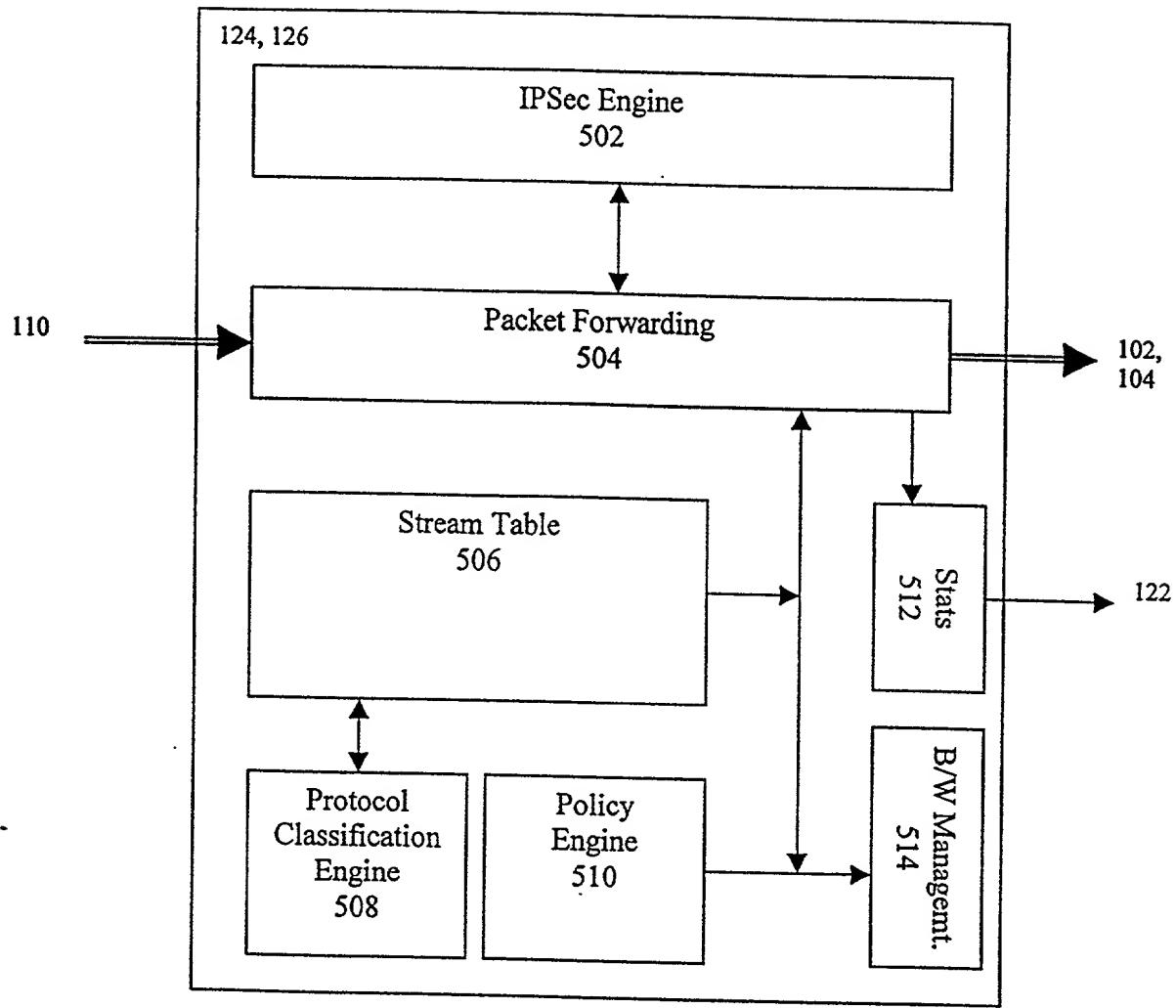


FIG. 18

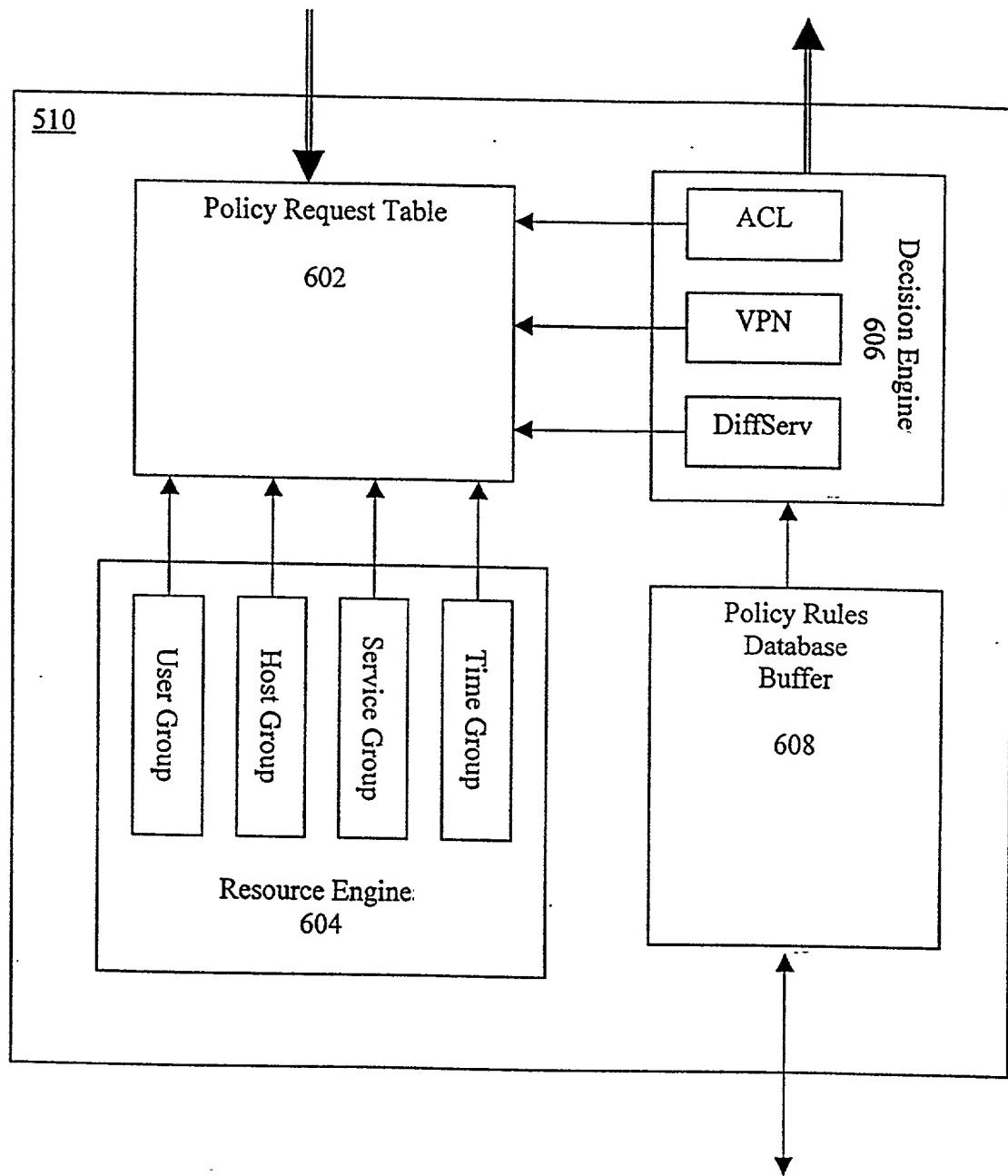


FIG. 19

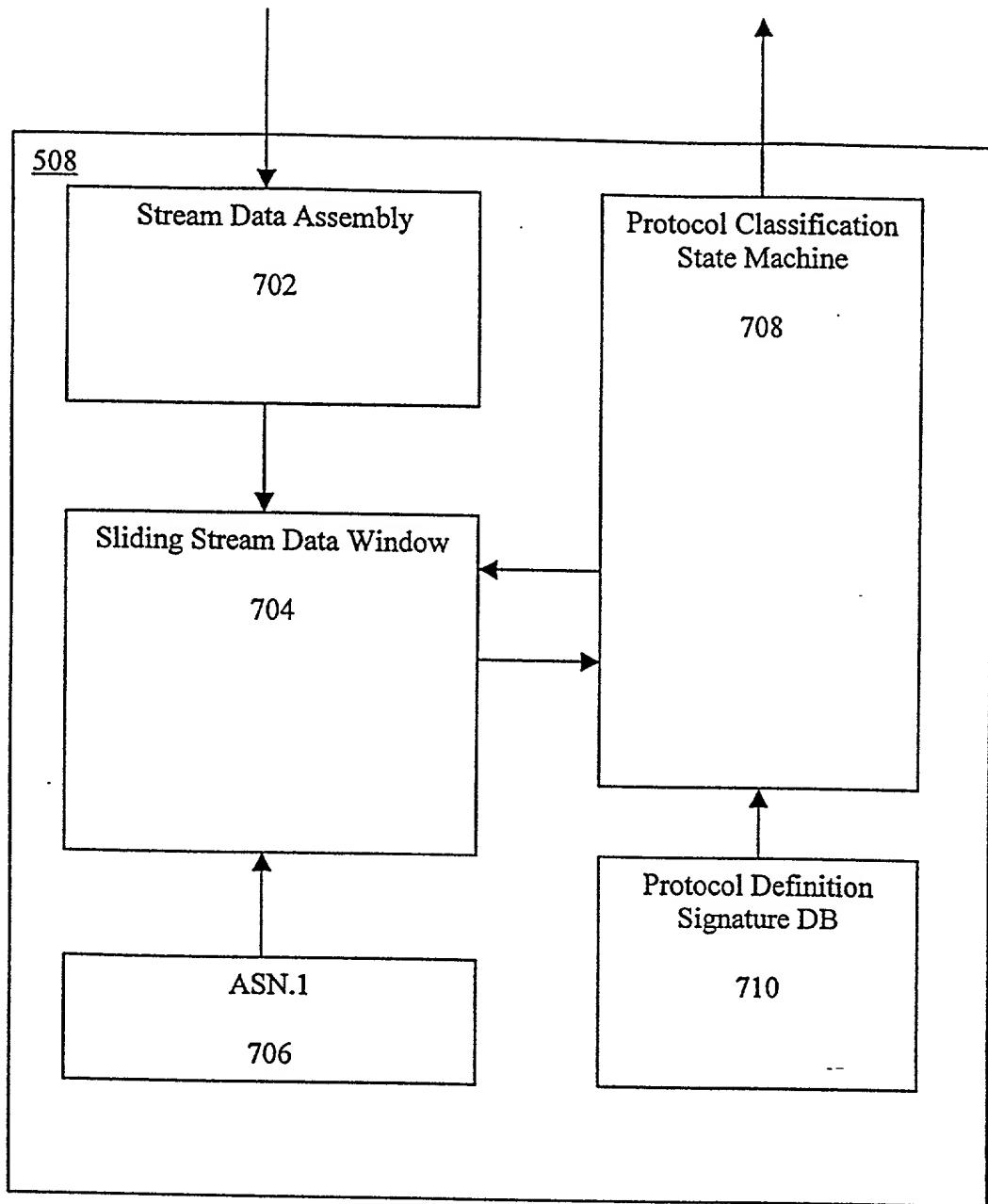


FIG. 20

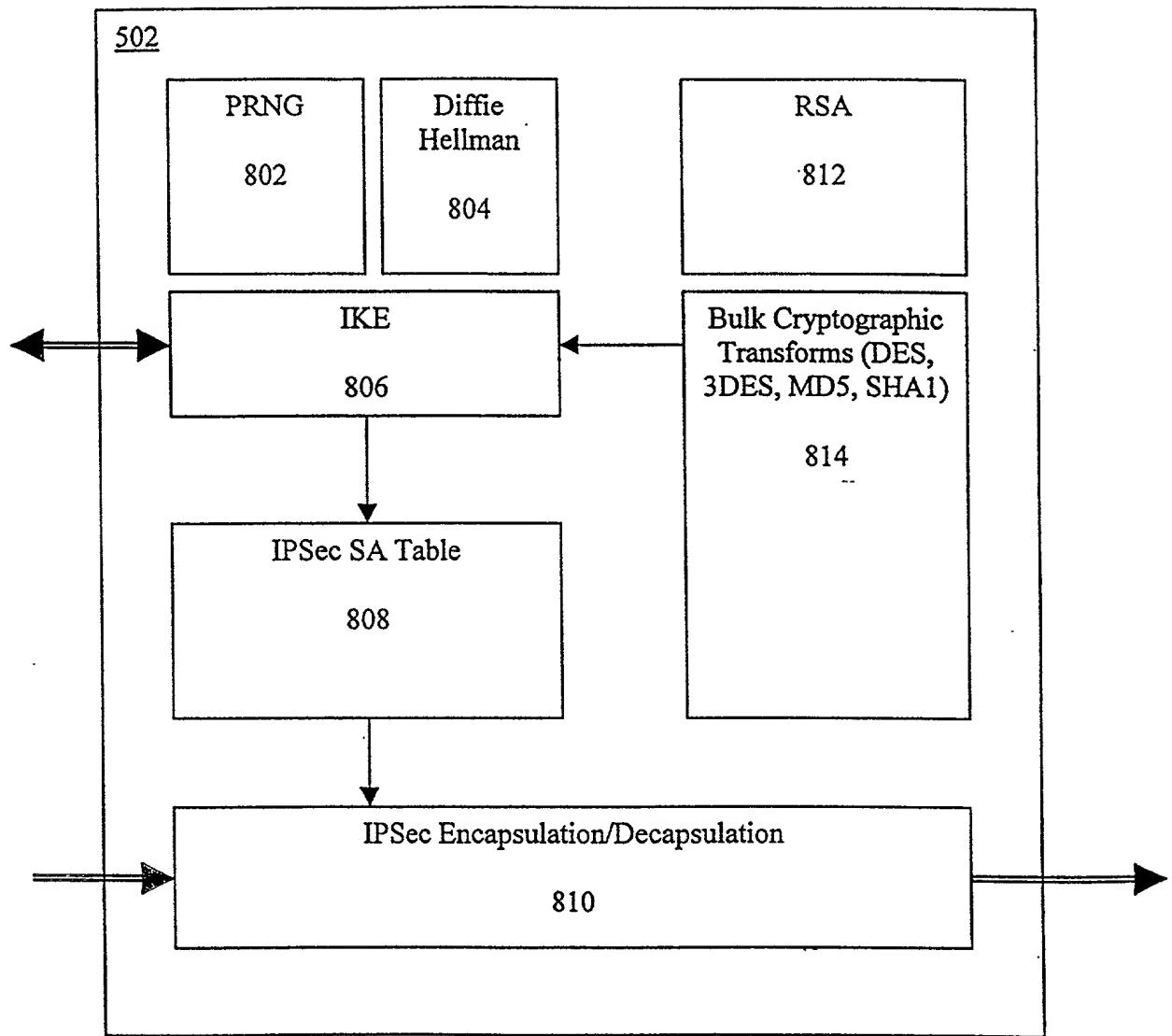


FIG. 21

Timestamp	Service	Source IP Address & Port	Destination IP Address & Port	User ID	Status	Operation	In Bytes	Out Bytes	Duration
820	821	822	823	824	825	826	827	828	829
820	821	822	823	824	825	826	827	828	829
820	821	822	823	824	825	826	827	828	829
820	821	822	823	824	825	826	827	828	829

FIG. 22

BRANCH B1
BRANCH B2
BRANCH B3
BRANCH B4

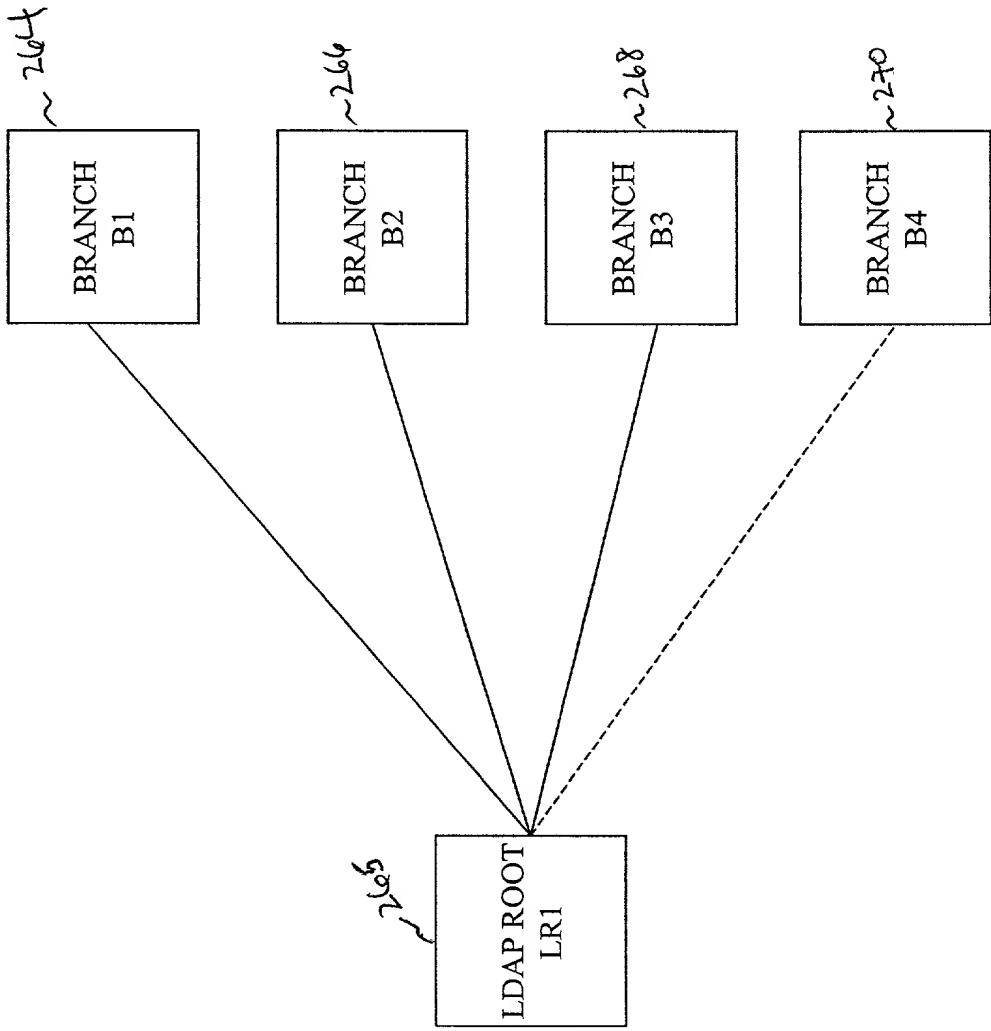


FIG. 23

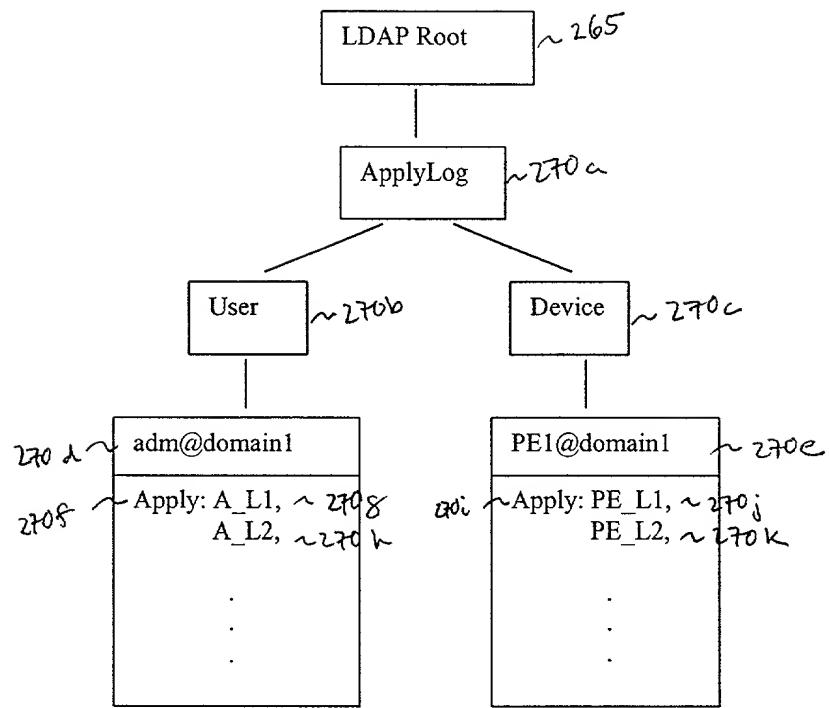
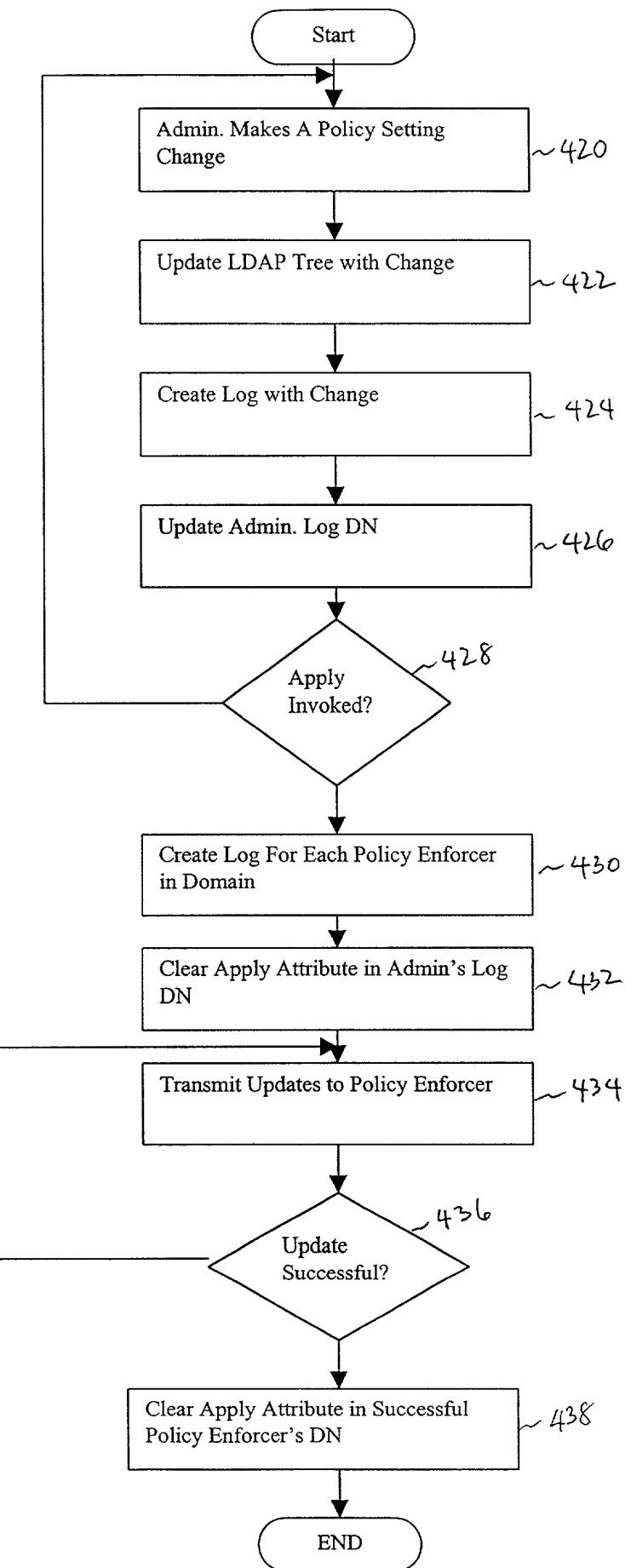


FIG. 24

FIG. 25



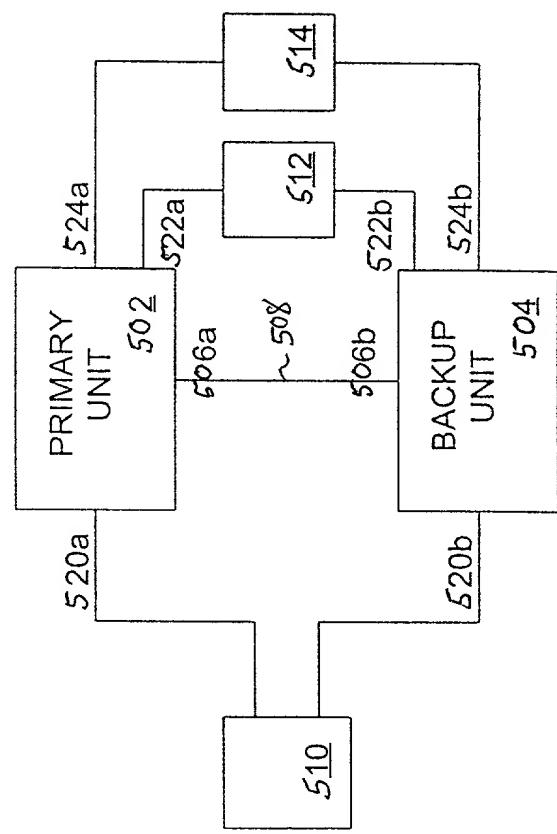


Figure 26

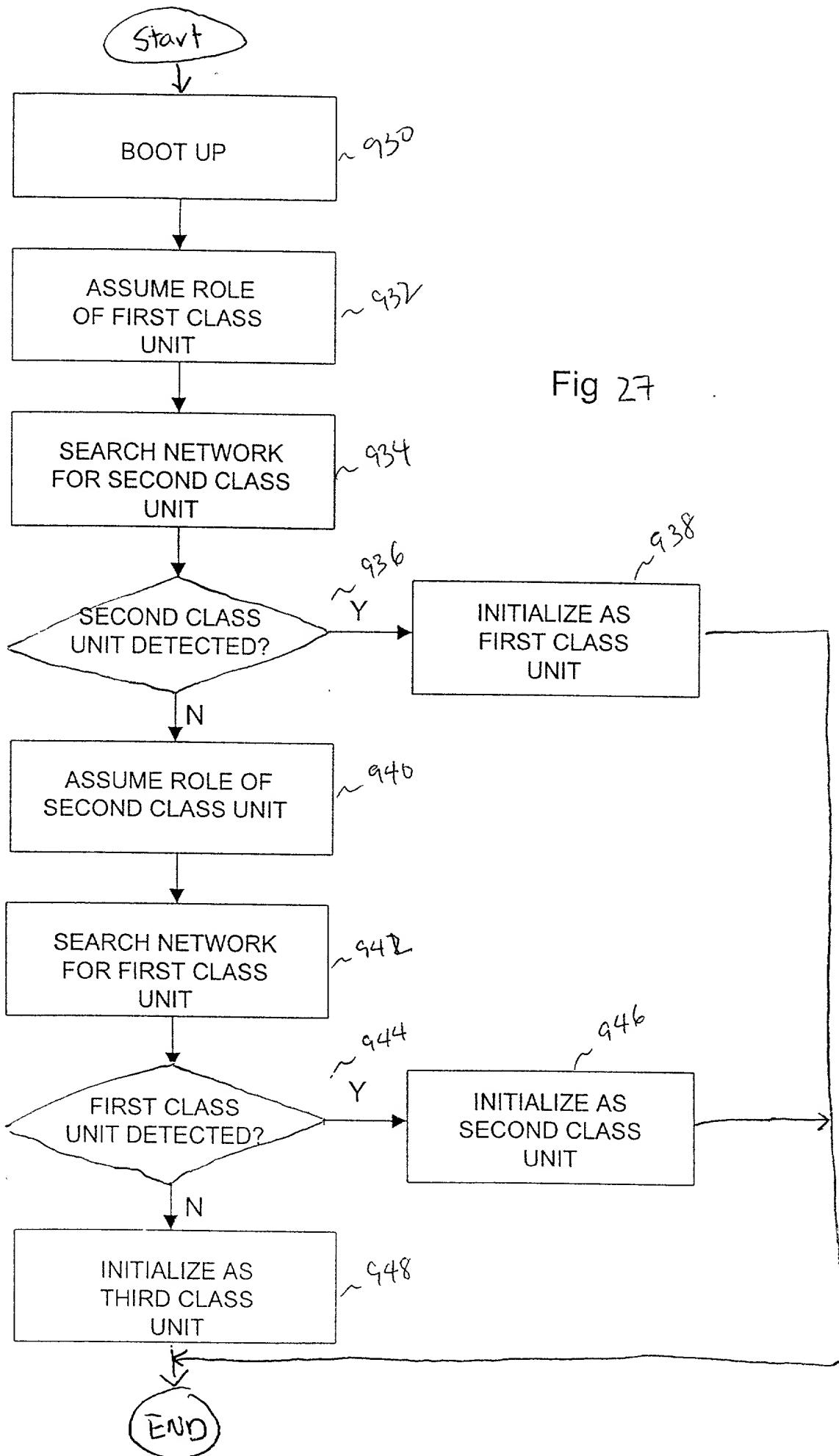


Fig 27

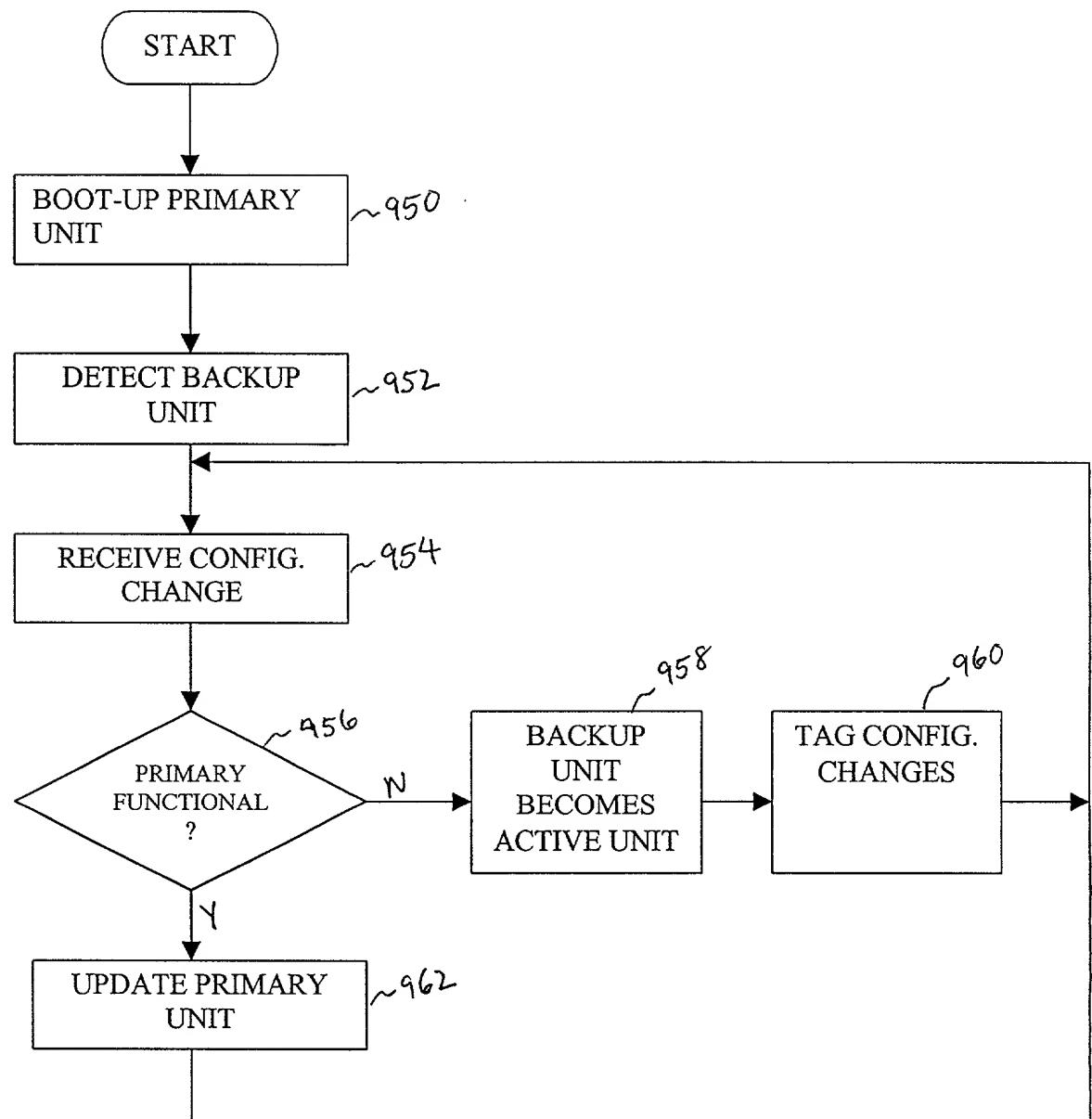


FIG. 28

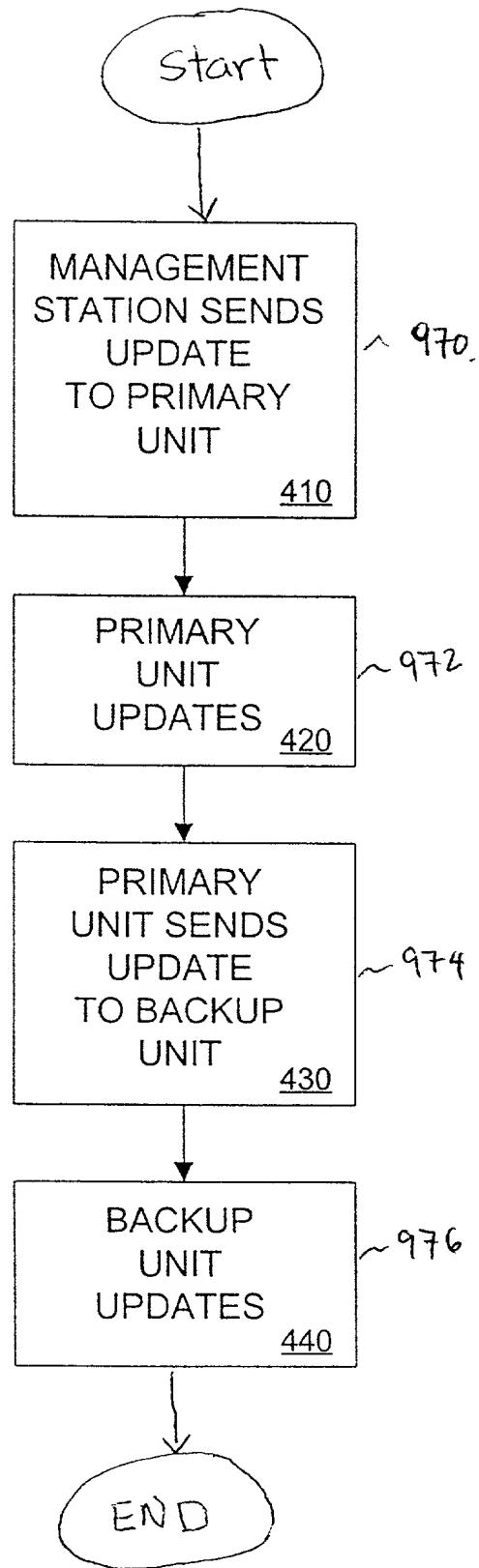


Fig 29

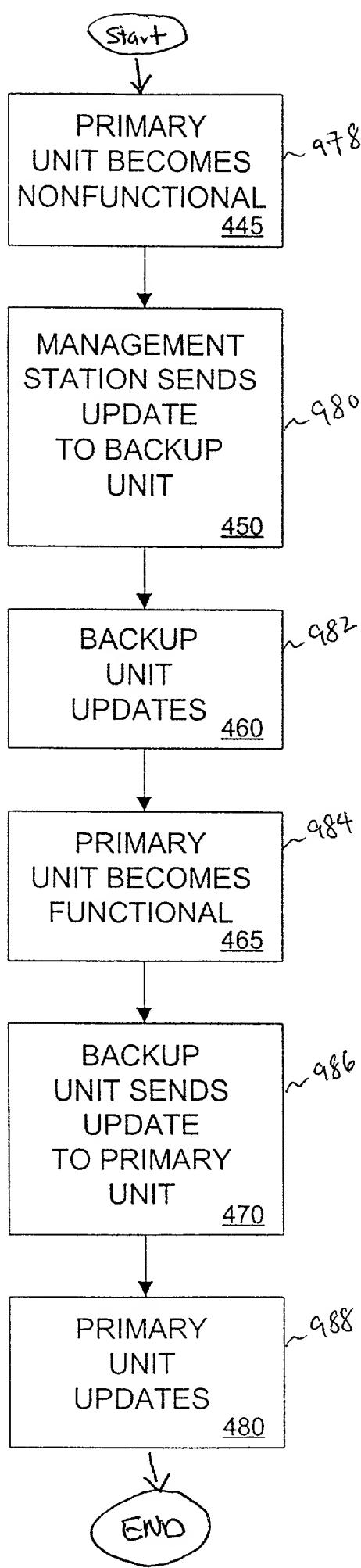


Fig. 30

**DECLARATION AND POWER OF ATTORNEY
FOR PATENT APPLICATIONS**

PATENT

Docket No. : 39299/JEC/X2

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name.

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled **OBJECT MODEL FOR NETWORK POLICY MANAGEMENT**, the specification of which is attached hereto unless the following is checked:

X was filed on June 12, 2000 as United States Application Number or PCT International Application Number (to be filled in when known) and was amended on (if applicable).

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose information which is material to patentability as defined in 37 CFR § 1.56.

I hereby claim foreign priority benefits under 35 U.S.C. § 119(a)-(d) or § 365(b) of the foreign application(s) for patent or inventor's certificate, or § 365(a) of any PCT International application which designated at least one country other than the United States, listed below and have also identified below, any foreign application for patent or inventor's certificate, or PCT International application having a filing date before that of the application on which priority is claimed.

Prior Foreign Application(s)

Application Number Country Filing Date (day/month/year) Priority Claimed

I hereby claim the benefit under 35 U.S.C. § 119(e) of any United States provisional application(s) listed below.

Application Number Filing Date

60/138,849	June 10, 1999
60,138,850	June 10, 1999
60/139,033	June 10, 1999
60/139,034	June 10, 1999
60/139,035	June 10, 1999
60/139,036	June 10, 1999
60/139,038	June 10, 1999
60/139,042	June 10, 1999
60/139,043	June 10, 1999
60/139,044	June 10, 1999
60/139,047	June 10, 1999
60/139,048	June 10, 1999
60/139,049	June 10, 1999
60/139,052	June 10, 1999
60/139,053	June 10, 1999
60/139,076	June 11, 1999

**DECLARATION AND POWER OF ATTORNEY
FOR PATENT APPLICATIONS**

Docket No. 39299/JEC/X2

I hereby claim the benefit under 35 U.S.C. § 120 of any United States application(s), or any PCT International application designating the United States, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the manner provided by the first paragraph of 35 U.S.C. § 112, I acknowledge the duty to disclose information which is material to patentability as defined in 37 CFR § 1.56 which became available between the filing date of the prior application and the national or PCT International filing date of this application:

<u>Application Number</u>	<u>Filing Date</u>	<u>Patented/Pending/Abandoned</u>
---------------------------	--------------------	-----------------------------------

POWER OF ATTORNEY: I hereby appoint the following attorneys and agents of the law firm CHRISTIE, PARKER & HALE, LLP to prosecute this application and any international application under the Patent Cooperation Treaty based on it and to transact all business in the U.S. Patent and Trademark Office connected with either of them in accordance with instructions from the assignee of the entire interest in this application; or from the first or sole inventor named below in the event the application is not assigned; or from in the event the power granted herein is for an application filed on behalf of a foreign attorney or agent.

R. W. Johnston	(17,968)	Gregory S. Lampert	(35,581)	Raymond R. Tabandeh	(43,945)
D. Bruce Prout	(20,958)	Grant T. Langton	(39,739)	Phuong-Quan Hoang	(41,839)
Hayden A. Carney	(22,653)	Constantine Marantidis	(39,759)	Kathy Mojibi	(41,409)
Richard J. Ward, Jr.	(24,187)	Marilyn R. Khorsandi	(45,744)	Cynthia A. Bonner	(44,548)
Russell R. Palmer, Jr.	(22,994)	Daniel R. Kimbell	(34,849)	Jun-Young E. Jeon	(43,693)
LeRoy T. Rahn	(20,356)	Craig A. Gelfound	(41,032)	Marc A. Karish	(44,816)
Richard D. Seibel	(22,134)	Syed A. Hasan	(41,057)	John F. O'Rourke	(38,985)
Walter G. Maxwell	(25,355)	Kathleen M. Olster	(42,052)	Richard J. Paciulan	(28,248)
William P. Christie	(29,371)	Daniel M. Cavanagh	(41,661)	Josephine E. Chang	(46,083)
David A. Dillard	(30,831)	Molly A. Holman	(40,022)	Frank L. Cire	(42,419)
Thomas J. Daly	(32,213)	Lucinda G. Auciello	(42,270)	Harold E. Wurst	(22,183)
Vincent G. Gioia	(19,959)	Norman E. Carte	(30,455)	Robert A. Green	(28,301)
Edward R. Schwartz	(31,135)	Joel A. Kauth	(41,886)	Derrick W. Reed	(40,138)
John D. Carpenter	(34,133)	Patrick Y. Ikehara	(42,681)	Stephen D. Burbach	(40,285)
David A. Plumley	(37,208)	Mark Garscia	(31,953)		
Wesley W. Monroe	(39,778)	Gary J. Nelson	(44,257)		

The authority under this Power of Attorney of each person named above shall automatically terminate and be revoked upon such person ceasing to be a member or associate of or of counsel to that law firm.

DIRECT TELEPHONE CALLS TO : Josephine E. Chang, 626/795-9900
SEND CORRESPONDENCE TO : CHRISTIE, PARKER & HALE, LLP
P.O. Box 7068, Pasadena, CA 91109-7068

I declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

**DECLARATION AND POWER OF ATTORNEY
FOR PATENT APPLICATIONS**

Docket No. 39299/JEC/X2

Full name of sole or first joint inventor Mahadevan Iyer	Inventor's signature	Date
Residence and Post Office Address 1075 Kildare Avenue, Sunnyvale, California 94087	Citizenship India	

Full name of second joint inventor Lavanya Apsani	Inventor's signature	Date
Residence and Post Office Address 3281 Falls Creek Drive, San Jose, California 95135	Citizenship India	

Full name of third joint inventor Pankaj Malviya	Inventor's signature	Date
Residence and Post Office Address 478 S. Fair Oaks Avenue, Sunnyvale, California 94086	Citizenship India	

DAA PAS255348.1-* 6/12/00 3 57 PM